

GAME THEORY - WEEK 1

FOCUS: ABSTRACT REASONING	HERMAN	TERMI
DISCIPLINING REASONING w/ MATHEMATICS	VITOR	
THE INTERNAL LOGIC OF MODELS	ADAM	WEI
DISCIPLINES OUR REASONING	KEVIN	(MILANCO)
"GAME THEORY DISCIPLINES THE CONCEPTS WE USE."	PATRICK	BONCA
	YUAN	GEOFF
	KENNEDY	TOM

ANALOGY w/ A LANGUAGE COURSE -> LOW GOAL IS TO KNOW VOCAB, ETC. SO NOT TO WRITE MODELS
 "IF YOU CAN'T WRITE IT, YOU DON'T KNOW IT."

* REFLECTION AT THE END OF PROBLEM SETS

GROWTH MINDSET VS. GRADES MINDSET
 -> PEOPLE WORKED WITH WEBSITES CONSULTED TABLES GENERATOR.COM - MATRICES + TABLES
 NO BLANKS ON HW -> TRY EVERYTHING TO GET A
 QUIZZES ON THE MONDAY AFTER PROBLEM SETS

KURAN CITES DAVID HACKETT FISHER ON HISTORIANS' FALLACIES

"POLITICAL REVOLUTIONS WILL CONTINUE TO SURPRISE US" B/C OF IRRATIONAL FALSIFICATION

SCIENTISTS: MODELS ARE 1) A PRECISE + ECONOMICAL STATEMENT OF A SET OF RELATIONSHIPS OR AN ACTOR (BIOLOGICAL, MECHANICAL, SOCIAL) SYSTEM

Why Models?

Game Theory - "How Instrumental
Rationality Manifests in Situations
of Strategic Interdependence"

In a strategic context, with multiple
people pursuing goal oriented behavior,
we can't just worry about the
probabilities (parameter reasoning), we
have to use models to guide
strategic reasoning

Model \neq Theory

↳ This is a class about models (of rationality)
Models are both concrete (specific
situations) and abstract (b/c they
are mathematically broad + all have
multiple applications).

Kuran - "Latent Revolutionary Bandwagon"
{0, 20, 30, 40, 50, 60, 70, 80, 100}
{0, 10, 20, 30, 40, 50, 60, 70, 80, 100}
TRAVOLUTION! No Revolution

Developments in Polisci often happen by
announcing from one discipline (legislative
bargaining) to another context (crisis bargaining)

Wed - Begin^{w/} Section 2 of Notes

GAME THEORY: TADDEUS CH 1

- 1) ACTIONS
- DECISION PROBLEMS: 2) OUTCOMES
- 3) PREFERENCES

PREFERENCE RELATIONS: $x \succeq y$: "x IS AT LEAST AS GOOD AS y"
 ↳ CAN BE EITHER:

STRICT PREFERENCE RELATION: $x \succ y$: "x IS STRICTLY BETTER"

INDIFFERENCE RELATION: $x \sim y$: "x AND y ARE EQUALLY GOOD"

COMPLETENESS AXIOM: ANY TWO OUTCOMES x, y CAN BE RANKED BY THE PREFERENCE RELATION SO THAT EITHER $x \succeq y$ OR $y \succeq x$.

- ↳ DOES NOT LET A PLAYER BE INDECISIVE
- ↳ ELIMINATES BURDEN'S JSS

TRANSITIVITY AXIOM: FOR ANY 3 OUTCOMES x, y, z , IF $x \succeq y$ AND $y \succeq z$, THEN $x \succeq z$

RATIONAL PREFERENCE RELATION: COMPLETE + TRANSITIVE

* NOTE CONSUMER PARADOX - RATIONAL PLAYERS IN A GROUP CAN FAIL TO REACH A DECISION
 "IMPOSING INDIVIDUAL RATIONALITY DOES NOT IMPLY GROUP RATIONALITY"

PROFIT/PAYOFF FUNCTION (RATIONAL PLAYERS ONLY)

EVERY ACTION $a \in A$ YIELDS A PROFIT $\pi(a)$

PAYOFF VALUES ARE ORDINATE \rightarrow NO ABSOLUTE VALUE REPRESENTED BY $u(\cdot)$ OR $v(\cdot)$

PLAYERS w/ RATIONAL PREFERENCES WILL MAXIMIZE THIS FUNCTION

PROPOSITION 1.1 \Rightarrow IF THE SET OF OUTCOMES X IS FINITE, THEN ANY RATIONAL PREFERENCE RELATION OVER X CAN BE REPRESENTED BY A PAYOFF FUNCTION.

DECISION TREES: TERMINAL NODE
DECISION NODE \circ
 \circ \longleftarrow \circ

RATIONAL CHOICE PARADIGM

\hookrightarrow ASSUMES FULL KNOWLEDGE OF POSSIBLE ACTIONS, OUTCOMES, AND PREFERENCES

PAYOFFS OVER ACTIONS: IF ACTION a LEADS TO OUTCOME $x(a)$, THEN THE PAYOFF FROM a IS GIVEN BY $v(a) = u(x(a))$, SO $v(a) =$ PAYOFF OF a

CONTINUOUS ACTION SPACE:

- 1) DEFINES PAYOFF FUNCTION
- 2) MAXIMIZE - TAKE 1ST DERIVATIVE + SET $= 0$

ETHICS + SOCIETY 13

TOLSON OFFICE CAN GIVE A LIST OF NAMES FOR SECTIONS

MORALITY IS ABOUT ACCOUNTABILITY - WE OUGHT TO BE ABLE TO HOLD OTHERS TO ACCOUNT.

GUIDE PARADIGMS STR. MUSK THEORY + MORAL THEORY
 ↳ WE HAVE USEFUL MORAL INTUITIONS, BUT (PACER BLOOM) THEY NEED TO BE DISCIPLINED BY THEORY

REVIEW IS- OUGHT DISTINCTION - ANY PROBLEMS?
 ↳ IS SCIENCE THEORY ONLY "IS"?

SEE HARRY SHAPIRO (!) ON IS-OUGHT

FOUR PSYCHOLOGY - USE RUSSELL ON THE GENERALIZABILITY OF IDEAS → SOMETHING WE WANT OTHERS TO WANT

JUSTICE - REVIEW MALCOLM LE'S TOKONOMY
 ↳ DISTRIBUTIVE, COMPLETIVE, ETC.

REVIEW SYLLABUS - HOW ARE THEY BEING EVALUATED?

DEONTIC LOGIC? SUPEREROGATORY? ("KINDS OF MORAL ACTS")

↳ PREFERENCE VS. OBLIGATION → IDEALS (AND NEG. IDEALS) ARE UNIVERSALIZABLE

HELP THEM SEE THE FOREST - PLAY THE VIDEOS HE SURFED
 WHAT IS A TOWN? - CONSIDERS!

GAME THEORY WEEK 1 READING

- 1) INSTRUMENTAL RATIONALITY - MEANS-END STRATEGIC INTERDEPENDENCE - DETERMINING ONE INDIVIDUAL'S OPTIMAL STRATEGY REQUIRES TAKING INTO ACCOUNT OTHERS' STRATEGIES.

DECISION THEORY BECOMES GAME THEORY ONCE YOU BRING IN STRATEGIC INTERDEPENDENCE.

- 2) A THEORY OF INSTRUMENTAL RATIONALITY MUST:

- 1) CAPTURE AN INTUITIVE IDEA OF CONSISTENT, GOAL-DIRECTED BEHAVIOR
- 2) SIMPLE - FOCUS ON ESSENTIAL ASPECTS
- 3) USEFUL - YIELDS NON-OBVIOUS INSIGHTS
- 4) RICH - INDICATES NEW APPROACHES
- 5) FLEXIBLE - WIDELY APPLICABLE

IT NEED NOT BE REALISTIC - PPL DON'T ALWAYS BE RATIONAL

- 3) GOAL-DIRECTEDNESS CAN BE MODELED AS PREFERENCES OVER A SET OF POSSIBLE OUTCOMES

1) OUTCOMES: $X = \{x_1, \dots, x_n\}$ NO OUTCOMES, IN ORDER
 (THE SET OF OUTCOMES IS FINITE)

2) PREFERENCES: REPRESENTED BY A TRAILING OF OUTCOMES: $R = \{(x, y) : x, y \in X\}$
 (ORDER MATTERS, NEED NOT INCLUDE ALL OUTCOMES)

INSTEAD OF $(x, y) \in R$, WE CAN WRITE

$x R y$ OR $x \succeq y$ - "x IS WEAKLY PREFERRED TO y"

$x \succ y$: "x IS STRICTLY PREFERRED TO y" ($x P y$)

$x \sim y$: "x IS INDIFFERENT TO y" ($x I y$)

ACTIONS: $A = \{a, b, c, \dots\}$ OR $A = \{a_1, a_2, \dots, a_k\}$
 (FINITE, FEASIBLE, CAN INCLUDE $(a+b)$ OR (ab))

4) CHOICE SET: $C(\sum_i x) \subseteq X$ (PERSON i 'S CHOICES)

\hookrightarrow SUBSET OF OUTCOMES i PREFERENCES AS MUCH OR MORE THAN OTHER OUTCOMES, DEFINED RELATIVE TO A PREF RELATION (\succeq_i) AND AN OUTCOME SET (X)

* (REVIEW EX. 4.11d)

$X \setminus \{w\}$ INCOMPLETE PREFS: THE CHOICE SET IS INCOMPLETE

WHenever preferences are 1) CYCLIC OR 2) INCOMPLETE

COMPLETENESS: A PREFERENCE RELATION IS COMPLETE

IFF FOR ALL $x, y \in X$, $x \succeq y$ OR $y \succeq x$

$\hookrightarrow \forall x, y \in X (x \succeq y \vee y \succeq x)$

TRANSITIVITY: A PREFERENCE RELATION IS TRANSITIVE

IFF FOR ALL $x, y, z \in X$, $x \succeq y \wedge y \succeq z$ MEANS $x \succeq z$

$\hookrightarrow \forall x, y, z \in X [(x \succeq y) \wedge (y \succeq z) \rightarrow x \succeq z]$

A RELATION THAT'S COMPLETE + TRANSITIVE IS CALLED AN ORDERING, A WEAK ORDER, OR A TOTAL PREFERENCE

* QUASI-TRANSITIVITY: $\forall x, y, z \in X [(x \succ y) \wedge (y \succ z) \rightarrow x \succ z]$

DEFINITIONS: $x \succeq y \leftrightarrow \neg (y \succ x)$

$x \sim y \leftrightarrow (\neg (y \succ x) \wedge \neg (x \succ y))$

ASYMMETRIC RELATIONS:

$\forall x, y \in X \neg (x \succ y \wedge y \succ x) \equiv \forall x, y \in X [x \succ y \rightarrow \neg (y \succ x)]$

NEGATIVE TRANSITIVITY:

$\forall x, y \in X [x \succ y \rightarrow \forall z \in X (x \succ z \vee z \succ y)]$

5) CONTINUOUS (INFINITELY DIVISIBLE) OUTCOME SETS:

\hookrightarrow COMPACT IFF CLOSED + BOUNDED

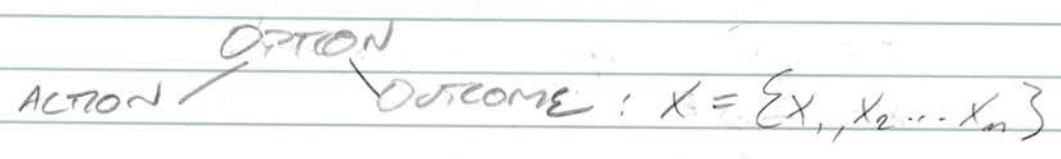
IF X IS INFINITE + NON-EMPTY, THEN IF X IS COMPACT

AND \succeq IS COMPLETE, TRANSITIVE + L.C., THEN $C(\sum_i x) \neq \emptyset$

GAME THEORY - WEEK 13 CLASS NOTES

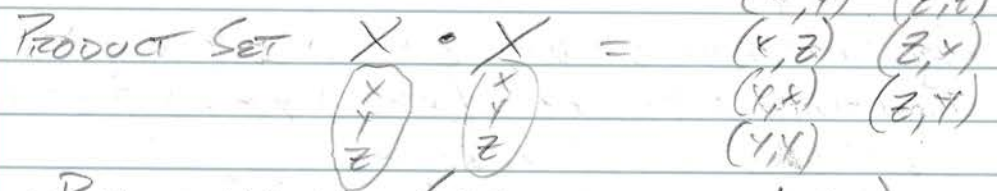
READ MEYERSON ARTICLE ON SCHEDULING

MEANS - END RATIONALITY - CONSISTENT, GOAL-DIRECTED BEHAVIOR



OUTCOMES PLUS PREFERENCES: ENDS
 ACTIONS: MEANS $A = \{a_1, \dots, a_n\}$

$R \rightarrow$ BINARY RELATIONS (ORDERED PAIRS)
 $(x, y) \neq (y, x)$ (x, x) (y, z)



$xRy = x \geq y$ (SUCCESSOR IN LATEX)



OUTCOMES: $X = \{x, y, z\}$
 PREFERENCES: $x \geq y, xRy$

$x \succ y \equiv x \geq y \wedge \neg(y \geq x)$
 $x \sim y \equiv x \geq y \wedge y \geq x$

AGENTS DO NOT HAVE PREFERENCES OVER ACTIONS, ONLY OVER OUTCOMES.

CHOICE SET: $C(\succeq, X)$

↳ ALWAYS DEFINED BY 1) CHOICE SET AND 2) OUTCOMES

$$C(\succeq, X) := \{x \in X : \forall y \in X [x \succeq y]\}$$

$\succeq = \begin{matrix} & & y & \\ & \nearrow & & \\ x & \longrightarrow & z & \end{matrix}$: $C = \{x, y\}$

$\begin{matrix} & & & \\ & & \searrow & \\ & \nearrow & & \\ x & \longrightarrow & z & \end{matrix}$: $C = \{\emptyset\}$ (EMPTY SET DOESN'T MEAN INDIFFERENCE)

NON-EMPTY CHOICE SETS CAN ARISE FROM

1) INCOMPLETENESS, AND 2) CYCLING

COMPLETENESS: $\forall x, y \in X (x \succeq y \vee y \succeq x)$

TRANSITIVITY: $\forall x, y, z \in X [(x \succeq y \wedge y \succeq z) \rightarrow x \succeq z]$

EDGE BTRW. EVERY NODE

IF X IS FINITE + NON-EMPTY AND \succeq IS COMPLETE AND TRANSITIVE, THEN THE CHOICE SET IS NOT EMPTY

PROOF: BY INDUCTION

ASSUME X IS FINITE + NON-EMPTY

BASE CASE: $X = \{x\}$: $C(\succeq, X) = \{x\}$

INDUCTION: $X' = \{x_1, \dots, x_n\}$: $C(\succeq, X') \neq \{\emptyset\}$

$$X = X' \cup \{a\}$$

$$= \{x_1, \dots, x_n, a\}$$

TWO CASES: $a \succeq x$ OR $x \succeq a$

BY TRANSITIVITY

$$C(\succeq, X) \neq \emptyset$$

(AT LEAST x)

ASK *

GAME THEORY 2a - NOTES

FUTURE CLASSES IN THE DEAN'S CONFERENCE ROOM
 ↳ CIRCULAR COURSEWORK, UP THE STAIRS TO THE LEFT

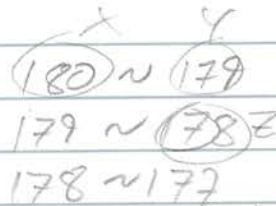
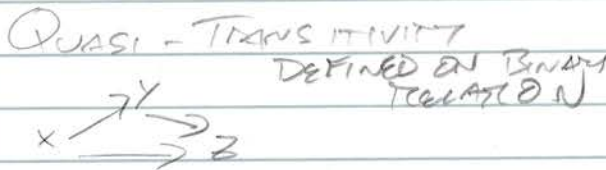
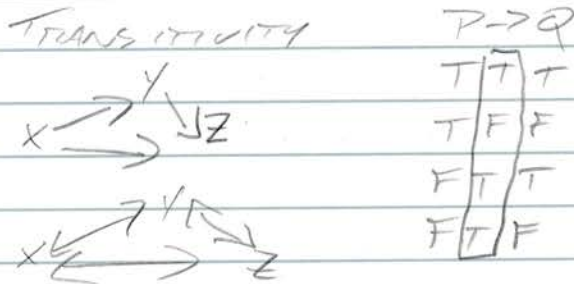
MBA122

BASE: $X = \{x\}$

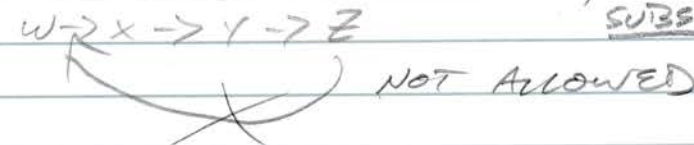
STEP 1: $X = \{x\} \cup \{a\}$ } $a \succeq x$
 $= \{x, a\}$ } $x \succ a$

CASE 1: $a \succeq x$ } IN BOTH CASES, CHOICE
 CASE 2: $x \succ a$ } SET IS FINITE + NON-EMPTY

COMPLETENESS + TRANSITIVITY ARE SUFFICIENT CONDITIONS FOR A NON-EMPTY CHOICE SET.



ACYCLICITY - NO CYCLES, DEFINED ON EVERY SUBSET OF OUTCOME SET.



Asx DW TO SEND PROOF OF 4.2

TRANSITIVITY $\xrightarrow{\text{IMPLIES}}$ QUASI-TRANS $\xrightarrow{\text{IMPLIES}}$ ACYCLICITY

4.2 ACYCLICITY IS BOTH NECESSARY & SUFFICIENT FOR CONSTRUCTING A NON-EMPTY CHOICE SET.

PROOF: SUPPOSE X IS FINITE & NON-EMPTY

SUPPOSE \succeq ON X IS COMPLETE

SHOW $C(\succeq, X) \neq \emptyset$ IFF \succeq IS ACYCLIC

$P \iff Q$

① IF $C \neq \emptyset$, THEN \succeq IS ACYCLIC

$P \implies Q$ - ASSUME ANTECEDENT!

ASSUME $C \neq \emptyset$

ASSUME \succeq NOT ACYCLIC

REDUCED TO ASSUMPTION

$X = \{w, x, y, z\}$ $C = \{w\}$

SAY $w \succ x \succ y \succ z$

IF $C = \{w\}$, THEN THERE CAN'T BE A CYCLE

THREE

MOST PROMINENT STRATEGIES FOR PROVING $P \implies Q$

ASSUME ANTECEDENT

(SEE VEULEMAN)

PROVE CONTRAPOSITIVE

AND ASSUME NEGATION OF CONSEQUENT

② IF \succeq IS ACYCLIC, THEN $C \neq \emptyset$

$w \succ x$

y

$w \succ y$

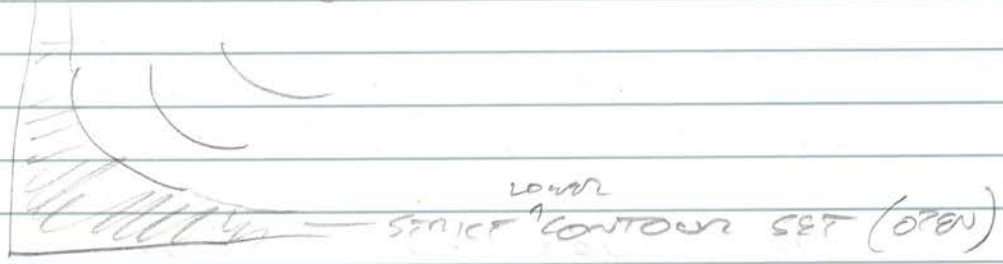
$y \succeq w$ OR $\neg(y \succeq w)$

EITHER $w \rightarrow x \rightarrow y$

OR $w \rightarrow x \leftarrow y$

5) CONTINUOUS CASE:
COMPACT: BOUNDED + CLOSED (ENDS ^{IN} IN SET)

*INDIFFERENCE CURVES



6) $u(x) \geq u(y)$ IFF $x \succeq y$
RELATION ON NUMBERS RELATION ON OUTCOMES

$$u(\emptyset) = \begin{cases} 7 & \emptyset = x \\ 3 & \emptyset = y \end{cases} \quad x \xrightarrow{y \succ x} z$$

$u(x) > u(z)$	$x > z$	} REPRESENTATION
$u(x) > u(y)$	$x > y$	
$u(y) = u(z)$	$y \sim z$	

WHAT MUST BE TRUE OF A PREFERENCE RELATION FOR IT TO BE REPRESENTED AS A UTILITY FUNCTION?

$\exists u: X \rightarrow \mathbb{R}$ THAT REPRESENTS \succeq IFF THAT PREFERENCE RELATION IS COMPLETE + TRANSITIVE.

(INTUITION: THE \succeq AND \prec RELATIONS ARE ALSO COMPLETE + TRANSITIVE)

PROOF: TAKE SOME FINITE X

FOR $X = X_i$, $U(x) = k - i$

(REPRESENTATION THEORY)



ORDINARY FUNCTIONS, IF TRANSFORMED,
CONVEY THE SAME INFORMATION.

SATIABLE PREFERENCES CAN BE MODELED BY

QUADRATIC $(t - t^*)^2$ OR $-|t - t^*|$

(DIFFERENTIABLE)

(NOT DIFFERENTIABLE)

(SPATIAL PREFERENCES)

10/10/18

GAME THEORY 213 NOTES

CARDINAL UTILITY FUNCTIONS

	COLD ^(.2)	MILD ^(.5)	HOT ^(.3)	OUTCOMES
BEACH	0	1	2	(BEACH, COLD)
MOVIE	2	0	1	(MOVIE, HOT)
HIKE	1	2	0	etc.

TAKE THE PROBABILITY-WEIGHTED SUM OF VALUES

$$\text{BEACH: } 0(.2) + 1(.5) + 2(.3) = 1.1$$

$$\text{MOVIE: } 2(.2) + 0(.5) + 1(.3) = .7$$

$$\text{HIKE: } 1(.2) + 2(.5) + 0(.3) = 1.2$$

THIS IS SUSCEPTIBLE TO A PROBLEM - TRANSFORMATION OF THE DATA WILL YIELD A DIFFERENT RESULT

SAME PREFERENCE RELATIONS, DIFFERENT NUMBERS

	COLD ^(.2)	MILD ^(.5)	HOT ^(.3)
BEACH	0	3	4
MOVIE	4	0	3
HIKE	3	4	0

$$\text{BEACH} = 2.7$$

$$\text{MOVIE} = 1.7$$

$$\text{HIKE} = 2.6$$

CARDINAL UTILITY FUNCTION: $u(x) - u(y) > u(x) - u(z)$

LOTTERIES: PROBABILITY DISTRIBUTIONS OVER OUTCOMES

BERNOULLI UTILITY FUNCTION: $u: X \rightarrow \mathbb{R}; \geq \text{or } <$

SET OF LOTTERIES: P ; INDIV. LOTTERIES: $\{p, q, r\}$

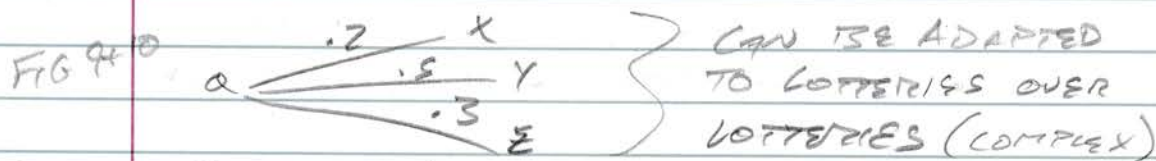
OR $\{p_1, p_2, p_3\}$

DEGENERATE LOTTERY: ONE OUTCOME HAS $p=1$

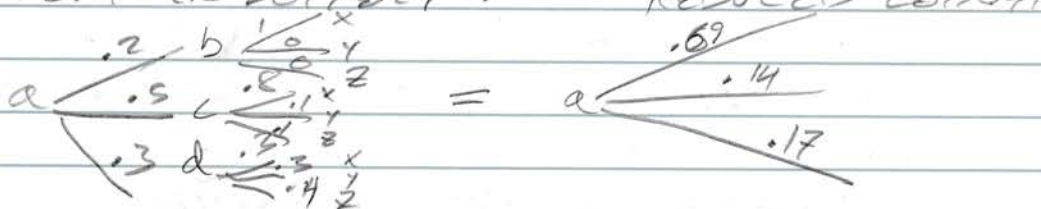
LOTTERIES ARE REPRESENTED AS VECTORS:

$\rightarrow (0.2, 0.5, 0.3; x, y, z)$

OR AS TREES:



COMPLEX LOTTERY:



CONSISTENT, GOAL-DIRECTED AGENTS CHOOSE THEIR MOST-PREFERRED LOTTERY.

(EXPECTED UTILITY MAXIMIZATION)

$P \succeq Q$ IF $p(x)u(x) + p(y)u(y) + p(z)u(z) \geq$ SAME FOR Q
 (LOTTERIES) (NUMBERS) $U(P) \geq U(Q)$

RUNNING FOR OFFICE

$u(\text{RUNNING} + \text{WINNING}) = 10 - 5 = 5$

$u(\text{RUNNING} + \text{LOSING}) = 0 - 5 = -5$

$u(\text{STATUS QUO}) = 0$



HOW MUCH DOES RUNNING FOR OFFICE HAVE TO RETURN?

PREFERENCES OVER LOTTERIES:

 \succsim ON \mathcal{P} (PREFERENCES OVER LOTTERIES)

PROOF OF PROPOSITION 9.1

VNM THEOREM

M

↓

M 9.1 \rightarrow [A PREFERENCE RELATION ON LOTTERIES
SATISFIES CONDITIONS (A, B, C, D) ONLY IF (IFF)]N THERE EXISTS A CARDINAL UTILITY RELATION
($U: X \rightarrow \mathbb{R}$) SUCH THAT FOR ALL LOTTERIES
IN THE SET OF LOTTERIES ($\forall p, q \in \mathcal{P}$) $p \succsim q$ IFF $\sum p(x) u(x) \geq \sum q(x) u(x)$

WHAT ARE AXIOMS A-D?

A: RATIONAL (COMPLETE, TRANSITIVE)

B: REDUCTION OF COMPOUND LOTTERIES

↳ PLAYERS CARE ABOUT OUTCOMES

C: CONTINUITY - $\alpha > \beta > \gamma$ $[\alpha p + (1-\alpha)r] \succsim q ; [\beta p + (1-\beta)r] \precsim q$ THE SET OF ALPHAS SUCH THAT p IS PREFERRED
TO q IS CLOSED

CONTINUITY RULES OUT LEXICOGRAPHIC PREFERENCES

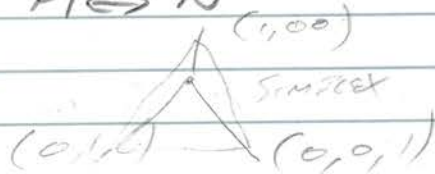
↳ EXTREME PREFERENCES \rightarrow NOT WILLING TO GAMBLED: INDEPENDENCE: PREFERENCES OVER TWO LOTTERIES
SHOULD REMAIN FIXED, EVEN IF WE MIX THE SAME
THIRD LOTTERY WITH BOTH OF THE FIRST TWO.

VON NEUMANN - MORGENTHAU AXIOMS (A-D)

TO PROVE (9.1) \rightarrow STRUCTURE: $M \stackrel{\text{IFF}}{\leftrightarrow} \mathbb{N}$ ASSUME THE AXIOMS: $A \succ B \succ C$

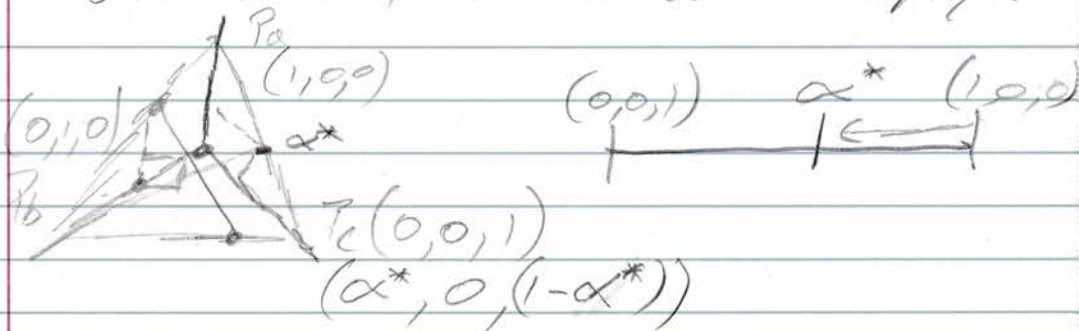
PROB. SIMPLEX IS THE FACE FACING US

+ PROBS ARE SORT TO 1



GAME THEORY 3A NOTES

10/15/18

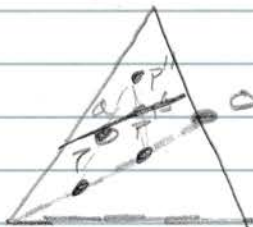
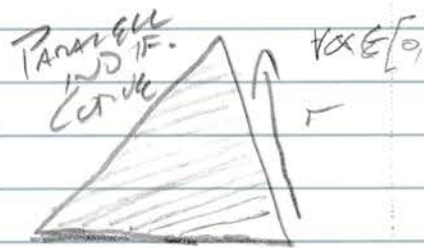
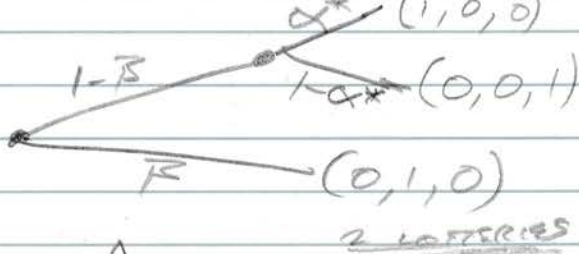


$$(0,1,0) \sim [\alpha^*(1,0,0) + (1-\alpha^*)(0,0,1)]$$

$$(0,1,0) \sim [\beta(0,1,0) + (1-\beta)(0,1,0)]$$

$$(0,1,0) \sim [\beta(0,1,0) + (1-\beta)[\alpha^*(1,0,0) + (1-\alpha^*)(0,0,1)]]$$

$$\beta(0,1,0) + (1-\beta)\alpha^*(1,0,0) + (1-\beta)(1-\alpha^*)(0,0,1)$$



$$\alpha^* \alpha \left[\frac{1}{2} P', \frac{1}{2} P'' \right] \quad (\text{VIA INDEPENDENCE})$$

$$\beta \left[\frac{1}{2} P', \frac{1}{2} P'' \right] \quad \text{AXIOM 1}$$

9.4.1 For any α , there's an indifference curve

$$U(P) = \alpha_P ; U(P) \geq U(Q) \text{ IFF } P \succeq Q$$

$$U(P) = U(Q) \text{ IFF } P \sim Q$$

$$U(P) \leq U(Q) \text{ IFF } P \preceq Q$$

$$U(P) = \sum_{x \in X} P(x) u(x)$$

10/15/18

GAME THEORY 3A NOTES CONT'D

We can now use a CARDINAL UTILITY FUNCTION TO REPRESENT PREFERENCES OVER LOTTERIES (AND THUS OVER OUTCOMES).

9.9 IF $U: X \rightarrow \mathbb{R}$ THAT REPRESENTS \succeq ON X , THEN $v(U(x)) = a \cdot U(x) + b$ ALSO REPRESENTS \succeq ON X . (a MUST BE POSITIVE, b ANYTHING)

↳ THE ONLY KIND OF TRANSFORMATION THAT PRESERVES CARDINAL UTILITY FUNCTIONS ARE LINEAR (AFFINE) TRANSFORMATIONS - THEY SHIFT THE SLOPE AND/OR INTERCEPT.

$$U'(p) = a \cdot U(p) + b$$

PREFERENCES ARE STILL PRIOR TO UTILITIES.

I DO NOT PREFER a OVER b BECAUSE IT HAS A HIGHER UTILITY, BUT RATHER a HAS A HIGHER UTILITY BECAUSE I PREFER IT.

WHY ON EARTH WOULD YOU ASSIGN A QUIZ QUESTION BASED ON HW Q'S MOST PPL GOT WRONG BEFORE REVISION?

STRATEGIC INTERDEPENDENCE

↳ INTRODUCING OTHER CHOOSING AGENTS
HOW DO WE MODEL THIS?

COOPERATIVE GAME: WHERE A "PRE-PLAY" STAGE PROVIDES AN OPPORTUNITY TO MAKE A (BINDING) AGREEMENTS ABOUT COOPERATION

NON-COOPERATIVE GAME: NO PRE-PLAY COLLUSION POSSIBLE, BUT COOPERATION IS STILL POSSIBLE, JUST NOT ENFORCEABLE
↳ OUR FOCUS WILL BE NON-COOPERATIVE GAMES (GAMES OF STRATEGY)

ELEMENTS: PLAYERS, ACTIONS, OUTCOMES, INFORMATION, TIMING

PRISONER'S DILEMMA: STANDARD EXPOSITION

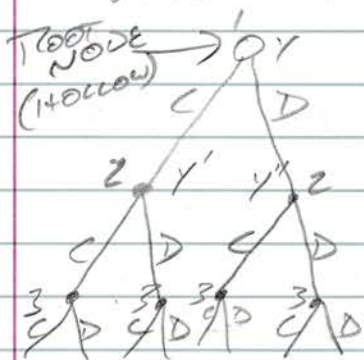
$I = \{1, 2, 3, \dots, n\}$ (PLAYERS)

u_i = UTILITY OF PLAYER i ; u_i = UTILITY OF P_i .

READ TO 3.5

GAME THEORY 3B - NOTES

SEQUENTIALLY



PRISONER'S DILEMMA

NODES: $Y = \{Y, Y', Y'' \dots\}$

Y' and Y'' HAVE NO RELATION BUT $Y > Y''$

↳ COMPLETE, TRANSITIVE ($Y > Y' > Y''$)

YOU CAN ONLY REACH SUBSIDIARY NODES VIA ONE ROUTE (UNIQUE TREE DECESSION)

ONE PLAYER AT EACH NODE

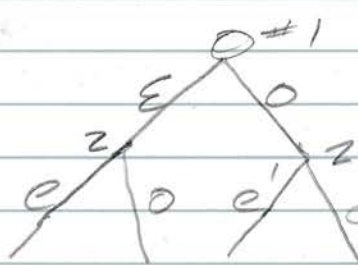
↳ FUNCTION: 1 TO 1

POWER SET: SET OF ALL SUBSETS

↳ SET: $\{A, B, C\}$ POWER SET: $\{\{A\}, \{B\}, \{C\}, \{A, B\}, \{A, C\}, \{B, C\}, \{A, B, C\}\}$

1) ASSIGN NODES TO PLAYERS ... AND \emptyset

2) ASSIGN ELEMENTS OF THE POWER SET TO NODES



(CANDIDATE ENTRY GAME)

ACTION SPACES:

$P_1: \{E, O\}$

$P_2: \{e, o, e', o'\}$

EACH OUTCOME IS UNIQUE: $(E, e) \neq (O, e')$

SIMULTANEOUS MOVE GAMES

COMPLETE VS. INCOMPLETE

PERFECT VS. IMPERFECT

CERTAIN VS. UNCERTAIN

INFORMATION SET: A SET OF NODES REPRESENTING

"WHAT A PLAYER KNOWS" → REPRESENTED BY ...

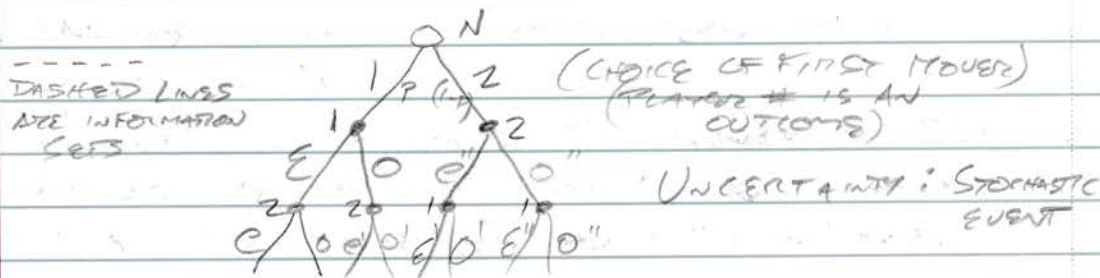
DOTTED LINES LINKING NODES. WHEN PLAYER

2 MOVES, THEY DON'T KNOW WHICH CONTINGENCY THEY'RE ACTING UNDER.

WHAT WE OBSERVE DEPENDS ON
WHAT HAPPENS COUNT OF ACTUALLY

IN THE CERTAIN CASE, ALL INFORMATION
SETS ARE SINGLETONS. IF THERE'S MORE
THAN ONE ELEMENT IN THE INFORMATION
SET, THEN THE CHOICE IS UNCERTAIN.

RANDOM CHANCE
NATURE: [^] MODELED AS A PLAYER IN THE GAME
↳ PROBABILITY DISTRIBUTION



COMMON KNOWLEDGE: THE PLAYERS HAVE A SHARED
UNDERSTANDING OF THE GAME.

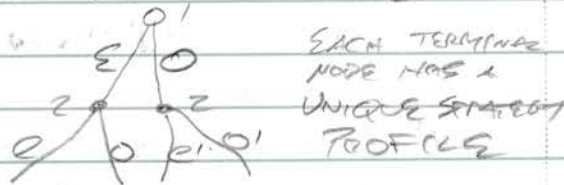
(THEY ALL KNOW THE FACTS, KNOW THAT EACH OTHER
PLAYER KNOWS THE FACTS, AND SO ON.)

PREFERENCES: UTILITY FUNCTIONS (DEFINED OVER
OUTCOMES)

BELIEFS: PROBABILITY DISTRIBUTIONS

STRATEGIES: A COMPLETE CONTINGENCY PLAN
↳ ASSIGNS AN ACTION TO EACH NODE

$S_1 = \{E, O\}$
 $S_2 = \{E, O, E', O'\}$



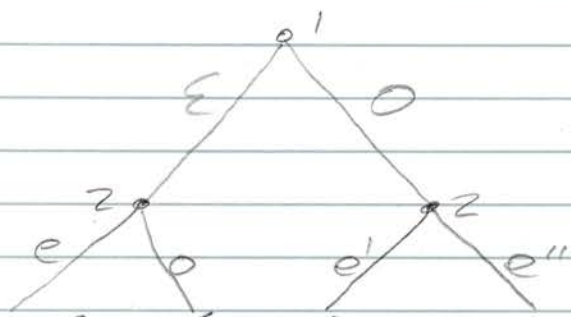
STRATEGY PROFILE:
CARTESIAN PRODUCT
OF STRATEGY SETS:

$S_1 \times S_2 = (E, ee'), (E, eo'), (E, oe'), (E, oo'), (O, ee'), (O, eo'), (O, oe'), (O, oo')$

GAME THEORY 4A - NOTES

EXPECTED UTILITY

CRUCIAL STRATEGY - COMPLETE, CONTINGENT PLAN FOR ALL PLAYERS
 COVERS BEST RESPONSE



IF A PLAYER HAS
 N NODES w/ M STRATEGIES,
 THEN THEY
 HAVE N^M STRATEGIES

$$S_1 = \{E, O\}$$

$$S_2 = \{ee', eo', oe', oo'\}$$

$$S_1 \times S_2 = S$$

	ee'	eo'	oe'	oo'
E	(E, ee')	(E, eo')	(E, oe')	(E, oo')
O	(O, ee')	(O, eo')	(O, oe')	(O, oo')

$$S = (S_1, S_2)$$

$$S = (S_1, S_2)$$

$$S_1 \times S_2 \times S_3 \times \dots \times S_n = N^M$$

Normal Form

$\frac{1}{2}$	ee'	eo'	oe'	oo'
E	1,1	1,1	2,0	2,0
O	0,2	0,0	0,2	0,0

EXTENSIVE FORM DOESN'T (NECESSARILY)
 INVOLVE TEMPORAL PRIORITY, BUT
 IS RATHER A SEQUENCE OF INFORMATION
 REVELATION.

STRATEGY PROFILES DEFINE UNIQUE
 PATHS THROUGH THE GAME TREE.

EXTENSIVE FORM REDUCES TO
NORMAL FORM (PLAYERS, STRATEGIES,
UTILITY FUNCTIONS).

STRATEGY PROFILES:

$$u_i(s_1, s_2) = \begin{cases} 1 & \text{IF } (E, ee') \text{ OR } (E, eo') \\ 0 & \text{IF - ALL THE REST (4)} \\ 2 & \text{IF } (E, oo') \text{ OR } (E, oo'') \end{cases}$$

EACH EXTENSIVE FORM HAS A UNIQUE NORMAL
FORM, BUT EACH NORMAL FORM CAN BE
REPRESENTED BY MULTIPLE EXTENSIVE FORMS.

THE SAME NORMAL FORM CAN REPRESENT DIFF. STRATEGIC SITUATIONS
COMMON PRACTICE: START WITH THE NORMAL FORM
TO Elicit INTUITIONS, THEN MOVE TO THE
EXTENSIVE (DYNAMIC) MODEL.

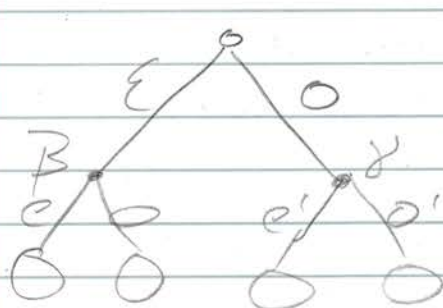
MIXED STRATEGIES: $S_1 + S_2$ (PREV. PAGE) ARE PURE
STRATEGIES. A MIXED STRATEGY IS A PURE STRATEGY
OVER A PROBABILITY DISTRIBUTION:

$P(ee') = .2$; $P(e, o') = .8$ (MUST SUM TO 1)
REPRESENTED BY SIGMA: $\sigma_1(ee') = .2$

BEHAVIOR STRATEGIES: ASSIGNS A DIFFERENT
PROB. DISTRIB. AT INFORMATION SETS.

(NORMAL
FORM
GAMES)

(



$$P(e, e') = .18$$

$$P(e, o') = .42$$

$$P(o, e') = .12$$

$$P(o, o') = .28$$

← could be any two numbers

IDENTIFYING WHICH STRATEGIES ARE RATIONAL

A RATIONAL PERSON CHOOSES THEIR MOST-PREFERRED LOTTERY GIVEN THE STRATEGIC CHOICES OF THE OTHER PLAYERS.

AN EQUILIBRIUM IS A STRATEGY PROFILE IN WHICH NO PLAYER WANTS TO UNILATERALLY DEVIATE (SPECIFIED BY STRATEGY PROFILE, NOT PAYOFFS)

NOTE: THE PRISONER'S DILEMMA EQUILIBRIUM IS NOT PARETO EFFICIENT - SEE GENERAL EQUILIBRIUM THEORY

(ME)

* THIS IS AN IMPLICIT CRITIQUE OF CAPITALISM

↳ ∃ SOME SITUATIONS WHERE DISTRIBUTED RATIONALITY FAILS TO REACH PARETO-OPTIMAL EQUILIBRIA.

SOLUTION CONCEPT - DOMINANT-STRATED REASONING

STRICT DOMINANT STRATEGY EQUILIBRIUM

↳ WHERE BOTH PLAYERS PLAY A STRICTLY DOMINANT STRATEGY

S_i STRICTLY DOMINATES S_i' IF
 IF $u_i(S_i, S_{-i}) > u_i(S_i', S_{-i})$
 FOR ALL $S_{-i} \in S_{-i}$ (NOT QUITE COMPLETE)
SEE 4.1-4.2

WEAK DOMINANCE: CHANGE $>$ (ABOVE)
 $\geq \rightarrow$ THEN

UTILITY OF A MIXED STRATEGY:

$$U_i(P, L) = P(2) + 1 - P(4) > 3$$

$$U_i(P, R) = P(6) + 1 - P(2) > 3$$

Solve for P

		L	R
(P)	U	2, 3	6, 2
	M	3, 1	3, 6
(1-P)	D	4, 0	2, 8

NO DOMINANT ^{PURE} STRATEGIES
 BUT THERE IS A DOMINANT
 MIXED STRATEGY

$$4 - 2P > 3; P < \frac{1}{2}$$

$$2 + 4P > 3; P > \frac{1}{4}$$



GAME THEORY HTB - NOTES

	L	C	R	Player 1
U	2,3	1,5	2,3	$u_1(D,L) > u_1(U,L) \quad S_1, S_2$
M	3,1	3,6	3,5	$u_1(D,C) > u_1(M,C) \quad S_1, S_2$
D	4,0	6,2	5,0	$u_1(D,R) > u_1(M,R) \quad S_1, S_2$

IF ALL 3 ABOVE ARE TRUE,
D STRICTLY DOMINATES U

IF D STRICTLY DOMINATES U,
BOTH U AND M, THEN
DIS A DOMINANT STRATEGY

STRICT DOMINANCE - MIXED STRATEGIES

	L	C	R
(P)	2,3	1,5	5,3
M	3,1	3,6	3,5
(1-P)	4,0	6,2	2,0

$$u_1(P,L) = P(2) + (1-P)(4) \quad 4 - 2P > 3$$

$$u_1(P,C) = P(1) + (1-P)(6) \quad 6 - 5P > 3$$

$$u_1(P,R) = P(5) + (1-P)(2) \quad 2 + 3P > 3$$

* → REVIEW INEQUALITY MANIPULATION $\frac{1}{2} < P < \frac{1}{2}$

DOMINANT STRATEGY (DOMINANCE) REASONING IS COARSE AND MINIMAL → IT'S A THRESHOLD CONDITION FOR RATIONALITY
(THIS SHOULD COME AFTER ITERATED DOMINANCE)

(Player 1 - Visual causal inference)

ITERATED DOMINANCE

	S_2
I_2 L C TR	TR DOMINATES C (Not L)
U 4,3 5,1 6,2	S_0 P_2 WILL NOT PLAY C
M 2,1 8,4 3,6	(SO WE WIPE OUT C)
D 3,0 9,6 2,8	NOW P_1 HAS A DOMINANT STRATEGY \rightarrow PLAY U
	(SO WIPE OUT M + D)
	AND WE END UP @ (4,3)

6.1 BEST RESPONSE (TO A SPECIFIC STRATEGY)

I_2 L C TR	S_1, S_2	S_1, S_2
U *1,1 -2,2 4,3*	$u_1(u, L) \geq u_1(M, L)$	S_1, S_2
M 0,3 3,1 5,4*	$u_1(u, L) \geq u_1(D, L)$	
D *1,5* *4,3* *6,2	U IS A BEST RESPONSE TO L	
	(SO IS D)	

BEST-RESPONSE ALGORITHM

\hookrightarrow PLAYER 1, COLUMN WISE (P_2 ROW WISE)

$S = \{D, L\}$ IS A CASE OF MUTUAL BEST RESPONSE.

EVEN THOUGH PLAYER 2 DOESN'T HAVE ANY

DOMINANT STRATEGY, C IS NEVER A B.R.

(SO WE ELIMINATE C AND M) (M WAS DOMINATED)

WHAT REMAINS IS THE SET OF RATIONALIZABLE STRATEGIES

GAME THEORY 413 - NOTES (CONT'D)

3-PLAYER Stag Hunt

		P = 2 1 3			
1	S	H		$u_2(S, S, S) \geq u_2(H, S, S)$	
	S	S	H	$3 \geq 2$	
2/3	$\begin{matrix} S \\ *3, 3, 3* \end{matrix}$	$\begin{matrix} H \\ 0, 0, 2 \end{matrix}$	$\begin{matrix} S \\ 2, 0, 0 \end{matrix}$	$\begin{matrix} H \\ *2, 0, 2* \end{matrix}$	$u_2(S, H, S) \geq u_2(H, H, S)$
	S	$\begin{matrix} 0, 2, 0 \\ 0, 2, 2 \end{matrix}$	$\begin{matrix} *2, 2, 0* \\ *1, 1, 1* \end{matrix}$		$0 \neq 2$

NASH EQUILIBRIUM

THE SOULTY RATIONAL STRATEGY PROFILE (WHERE EACH PLAYER IS PLAYING A BEST RESPONSE). (ABOVE IN BOXES)
 "A STRATEGY PROFILE OF MUTUAL BEST RESPONSES"

BEST RESPONSE CORRESPONDENCE

$$BR_1(S_{-1}) = \begin{cases} S & \text{if } (S, S) \\ H & \text{if } (S, H) (H, S) \\ H & \text{if } (H, H) \end{cases} \text{ OR } \begin{cases} S & \text{if } (SS) \\ H & \text{otherwise} \end{cases}$$

	U	C	D	
U	$5, 1$	$1, 4$	$1, 0$	$BR_2(S_2) \begin{cases} \{U\} & \text{if } S_2 = U \\ \{M, D\} & \text{if } S_2 = M \\ \{M\} & \text{if } S_2 = D \end{cases}$
M	$3, 2$	$*4, 0*$	$*3, 5*$	
D	$4, 3$	$*4, 4*$	$0, 4$	

NASH EQUILIBRIA ARE ROBUST TO ITERATIVE REASONING AND DON'T REQUIRE THE COMMON KNOWLEDGE ASSUMPTION (IT IS REQUIRED FOR ITERATED DOMINANCE).

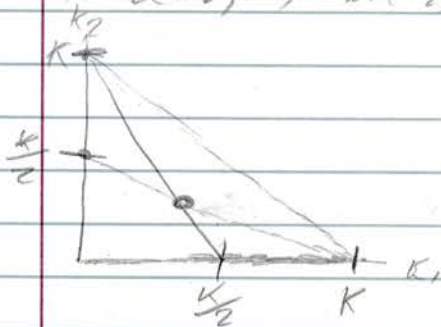
GAME THEORY SA - NOTES 10/29/12

CONDITIONS FOR WRITING A UTILITY FUNCTION
(PREFER TRAIL SUMMATING TEST RESPONSES)

2 FIRMS THAT CONSUME CLEAN AIR

$$u_1(k_1, k_2) = \ln(k_1) + \ln(K - k_1 - k_2)$$

$$u_2(k_1, k_2) = \ln(k_2) + \ln(K - k_1 - k_2)$$



(CORRESP.)

BEST RESPONSE FUNCTION?

↳ FIRST DERIVATIVE MAXIMIZES

$$\frac{\partial u_1}{\partial k_1} = \frac{1}{k_1} - \frac{1}{K - k_1 - k_2}$$

$$k_1^*(k_2) = \frac{K - k_2}{2} \quad (\text{BEST RESPONSE CORRESPONDENCE})$$

2 EQUATIONS,

2 UNKNOWN

$$k_2^*(k_1) = \frac{K - k_1}{2}$$

↳ SOLVE FOR $k_1^* + k_2^*$; $k_1^* + k_2^* = \frac{K}{3}$

- PROCEDURE:
- 1) FIND BEST RESPONSE
 - 2) TAKE 1ST DERIVATIVE
 - 3) SOLVE FOR k_1, k_2 (MAXIMIZE)
 - 4) SOLVE SYSTEM (FIND INTERSECTION)

MIXED STRATEGY NASH EQUILIBRIA

H T

H $\begin{matrix} * & 1 & -1 \\ * & -1 & * \end{matrix}$ NO PURE STRATEGY

T $\begin{matrix} -1 & * & * \\ * & -1 & -1 \end{matrix}$ $P_2: q \text{ ON H, } (1-q) \text{ ON T}$

GIVEN THIS, FIND A q FOR T , SUCH THAT T IS INDIFFERENT IS FOR H & T

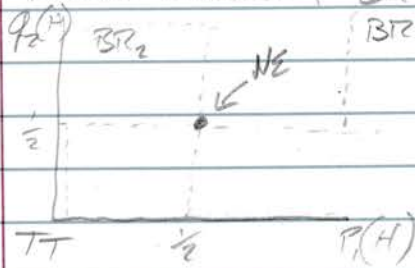
$$u_1(H, q) = q - 1 + q = 2q - 1 \quad \left. \begin{array}{l} 2q - 1 \geq 1 - 2q \\ q \geq \frac{1}{2} \end{array} \right\}$$

$$u_1(T, q) = -q + (1 - q) = 1 - 2q$$

u_2

u_2

MIXED STRATEGY EQUILIBRIUM: $P = \frac{1}{2}, Q = \frac{1}{2}$

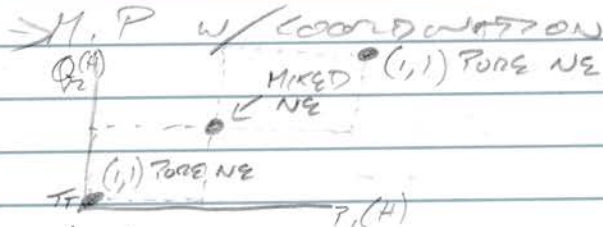


WHEN $P = \frac{1}{2}$, P_2 IS INDIFFERENT BETWEEN ANY MIX OF $Q + 1 - Q$ (AND VICE VERSA)

↳ THESE ARE THE VERTICAL + HORIZONTAL LINES



	H	T
H	$\frac{1}{2}, \frac{1}{2}$	0, 0
T	0, 0	$\frac{1}{2}, \frac{1}{2}$



$$U_1(P, Q) = \frac{1}{2} \left[\frac{1}{2}(1) + \frac{1}{2}(0) \right] + \frac{1}{2} \left[\frac{1}{2}(0) + \frac{1}{2}(1) \right] = \frac{1}{4} + \frac{1}{4} = \frac{1}{2}$$

IN THIS CASE, THE PURE STRATEGY EQ. ARE PATTERO-SUPERIOR TO THE MIXED STRATEGY EQ.

PLAYERS ONLY PLAY MIXED STRATEGY WHEN THEY ARE INDIFFERENT AMONG THEIR PURE STRATEGIES.

BOTH PLAYERS NEED NOT PLAY MIXED STRATEGIES
↳ ONLY ONE CAN, AND IT'S STILL A M.S.S.

ROCK PAPER SCISSORS

	TR	P	S
TR	0, 0	-1, 1	1, 1
P	1, -1	0, 0	-1, 1
S	-1, 1	1, -1	0, 0

ASSUME
- PLAYER 2 RULES OUT TR
THEN P₁ WILL NEVER PLAY P
AS A RESULT, P₂ WILL EXCLUDE TR

PLAYERS NEVER PLAY A MIX ✓ DOMINATED STRATEGY

THE UPSHOT OF THIS EXAMPLE IS THAT PLAYERS WILL ALWAYS BE FORCED TO MIX EQUALLY AMONG ALL THE NONDOMINATED STRATEGIES AVAILABLE TO THEM.

↳ AND NON-BEST RESPONSES

SYMMETRIC MIXED STRATEGY EQUILIBRIUM
3-PLAYER MIXED STRATEGIES (SEE HW NOTES)

	1	2	3	
	E	N		
1	E	N	E	N
E	$(\frac{1}{3}, \frac{1}{3}, \frac{1}{3})$	$(0, 0, 0)$	$(0, 0, 0)$	$(0, 1, 0)$
N	$(0, 0, 0)$	$(1, 0, 0)$	$(0, 0, 1)$	$(\frac{2}{3}, \frac{2}{3}, \frac{2}{3})$

$$P_2 = P_3 = P$$

$$P_1: Pr(E) = P^2$$

$$Pr(E) = P(1-P)$$

$$Pr(N) = P(1-P)$$

$$Pr(NN) = (1-P)^2$$

$$u_1(E, P_2, P_3) = P^2(-\frac{1}{3}) + P(1-P)(0) + P(1-P)(0) + (1-P)^2(-1)$$

$$= 1 - 2P + \frac{2}{3}P^2$$

$$u_1(N, P_2, P_3) = \frac{4}{3}P + \frac{2}{3}P^2$$

$$1 - 2P + \frac{2}{3}P^2 = \frac{4}{3}P + \frac{2}{3}P^2$$

SET EQUAL?
SEE NOTES

QUADRATIC EQUATION: YIELDS 2 SOLUTIONS

↳ USE THE ONE (S) BTW. 0 + 1

(COULD BE BOTH, 1, OR NEITHER)

10/51/18

HOW TO INTERPRET MIXED-STRATEGY EQUILIBRIA

↳ REVELATION CONDITION: IF IT WERE REVEALED THAT THE OTHER PLAYER HAS A PARTICULAR

NASH'S FUNDAMENTAL THEOREM - EVERY FINITE GAME HAS AT LEAST ONE MIXED STRATEGY NE.

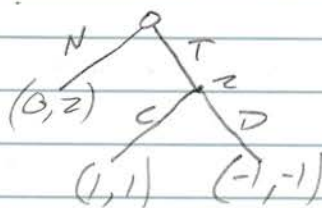
STRATEGIC RATIONALITY - GOAL-DIRECTED AGENTS CHOOSE A UTILITY-MAXIMIZING STRATEGY GIVEN WHAT OTHER AGENTS ARE DOING.

SEE NOTES

GAME THEORY 5B - NOTES

10/31/18

DYNAMIC GAMES w/ COMPLETE INFORMATION



	C	D
N	0, 2*	0, 2*
T	1, 1*	-1, -1

SUSTAINED BY
A NON-CREDIBLE
THREAT

THE NASH EQUILIBRIUM PREDICTS THAT ND IS A BEST RESPONSE \rightarrow BUT IF P_2 WERE CHOOSING, THEY WOULD PICK C.

NASH EQUILIBRIUM DELIVERS COUNTERINTUITIVE RESULTS IN DYNAMIC SITUATIONS.

PATH OF PLAY \rightarrow THE SET OF NODES REACHED WITH POSITIVE PROBABILITY IN A PARTICULAR STRATEGY PROFILE. IF THAT STRATEGY PROFILE IS IN EQUILIBRIUM, WE CALL IT THE EQUILIBRIUM PATH OF PLAY. ABOVE, N, D IS ONLY A NASH EQUILIBRIUM BECAUSE IT KNOWS WHAT'S HAPPENING OFF THE PATH OF PLAY.

THE PATH OF PLAY ENDS WHEN THERE'S A DISCONTINUITY IN THE PATH AS WRITTEN ON A TREE. THIS DISTINGUISHES BETWEEN THE ACTUAL PATH (ROOT TO END) AND THE COUNTERFACTUAL PATH.

3.1 SUBGAME PERFECTION - SEQUENTIAL RATIONALITY
↳ A STRATEGY IS SEQUENTIALLY RATIONAL IF IT IS RATIONAL AT EVERY NODE IN THE PATH OF PLAY.

FOR A PLAYER'S STRATEGY TO BE ^{SEQUENTIALLY} RATIONAL, THEY MUST PLAY A BEST RESPONSE AT EACH NODE.

CONTINUATION GAME VS. CONTINUATION PAYOFF → A NODE WHERE A PLAYER IS CHOOSING BETWEEN A CERTAIN PAYOFF AND A FURTHER SUBGAME. (BACKWARDS INDUCTION)

A SEQUENTIALLY RATIONAL STRATEGY IS ONE THAT CONSISTENTLY CHOOSES THE HIGHEST CONTINUATION PAYOFF. A PROFILE IS SEQUENTIALLY RATIONAL IF IT REQUIRES PLAYERS TO PLAY A BR @ EACH NODE.

REFINEMENT: PLAYERS SHOULD PLAY A NASH EQ. AT EACH CONTINUATION GAME. SEQUENTIALLY RATIONAL NEs ARE A SUBSET OF THE FULL SET OF NEs.

BACKWARDS INDUCTION - BEGIN AT THE FINAL SET OF NODES, & FIND SEQUENTIALLY RATIONAL RESPONSES AT EACH. THEN TRACE THE PATH OF PLAY.

1 / 1

SEQUENTIALLY RATIONAL STRATEGY PROFILES
MUST ACCOUNT FOR COUNTERFACTUALS.

BECAUSE WHAT HAPPENS OFF THE PATH
OF PLAY DETERMINES THE PATH OF PLAY

FOR EVERY FINITE GAME OF PERFECT INFORMATION
THERE EXISTS ^(AT LEAST ONE) A BACKWARDS INDUCTION SOLUTION
IF NO NODES HAVE EQUAL PAYOFFS, THAT
SOLUTION IS UNIQUE.

PROPER SUBGAME - STARTS AT A SINGLETON
NODE + INCLUDES ^{ALL} SUCCESSOR NODES, AND AVOID
CUTTING ANY INFORMATION SETS. THE ENTIRE
GAME COUNTS AS A PROPER SUBGAME <sup>(DISTINCT FROM PROPER
SUBSET)</sup>

PLAYERS PREFER BEST RESPONSE IN EACH
PROPER SUBGAME.

GAME THEORY WEEK 6A

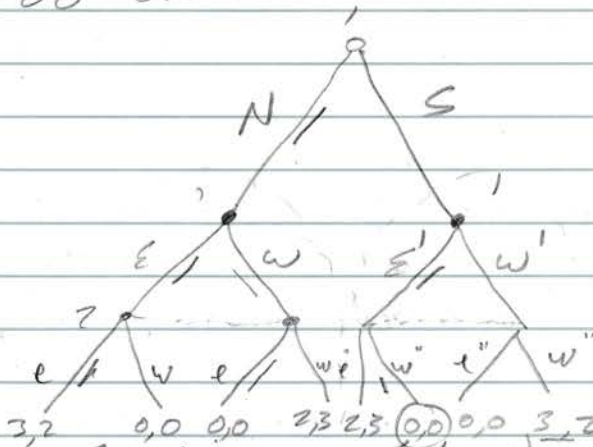
11 / 5 / 18

SUBGAME PERFECT EQUILIBRIUM

- ↳ PLAYERS PLAY AN NE IN EACH PROPER SUBGAME
- ↳ STRATEGY PROFILE, RESTRICTED TO EACH SUBGAME, IS A NASH EQUILIBRIUM.

$\sigma = (NE^1, we^1)$ (RESTRICTION TO SUBGAME)

PROCEDURE: GO DOWN TREE, TAKE THE LOWEST SUBGAME, FIND ITS NASH EQ., PRESERVE AS WE GO UP.



(NE^1, e^1)

(NE^1, ew^1)

NE: (MOST IN CLUDE) (W, w) (MODEL NOUANTE FORM TO FIND)

	ee'	ew'	we'	ww'
NSE	(3,2)*	(3,2)*	0,0	0,0
NEW	(3,2)*	(3,2)*	0,0	0,0
NWE	0,0	0,0	(2,3)*	2,3*
NWW	0,0	0,0	(2,3)*	2,3*
SEE'	2,3*	0,0	(2,3)*	0,0
SEW'	0,0	(3,2)	0,0	(3,2)
SWE'	2,3*	0,0	(2,3)	0,0
SWW'	0,0	(3,2)*	0,0	(3,2)

CIRCLED EQUILIBRIA FIT THE NE REQUIREMENTS ABOVE; THE OTHERS DON'T

SPESENE

DEFINITION: CONTINUATION VALUE?

SINCE DEVIATION PRINCIPLE - NASH EQUILIBRIUM
 FOR SUBGAME PERFECT IF NEITHER PLAYER
 HAS AN INCENTIVE TO DEVIATE FROM THE
 PATH OF PLAY. (ONE DEVIATION IS ENOUGH
 TO TAKE OUT SUBGAME PERFECTION).
 ↳ IMPORTANT TO HOLD EVERYTHING ELSE FIXED.

SUBGAME PERFECT ONLY IF NO PLAYER HAS
 A SINGLE PROFITABLE DEVIATION.
 USEFUL FOR REPEATED GAMES (EVEN INFINITE GAMES)

MULTI-STAGE GAMES - DYNAMIC INTERACTIONS
 ↳ PAYOFFS COME AS A SEQUENCE
 SECF-CONTAINED INTERACTIONS WITH PAYOFFS
 UNFOLDING OVER TIME.

L = LEAVE ALONE
 P = PUNISH

L = LEAVE ALONE
 P = PUNISH

SEE PROP. 5.1

	L	D	L	P	PERCEPTIONS @ LATER STAGES
C	4,4	-1,5	0,0	-3,-1	INFLUENCE / CONDITION
D	5,1	1,1	-1,3	-2,-2	BEHAVIOR @ EARLY STAGES

PLAYERS WILL MAXIMIZE UTILITY (GIVEN OTHERS' STRATEGIES)
 OVER ALL STAGES OF THE GAME.

PLAYERS WILL DISCOUNT BENEFITS OVER TIME

$$\begin{aligned}
 & x^1, x^2, x^3, \dots, x^T ; u(x^T, \delta) = \\
 & u(x^T, \delta) = u_1(x^1) + \delta u_1(x^2) + \delta^2 u_1(x^3) + \dots \\
 & + \dots + \delta^{T-1} u_1(x^T) \\
 & (x + \delta x + \delta^2 x + \delta^3 x + \dots + \delta^{T-1} x) = \sum_{t=1}^T \delta^{t-1} x (1-\delta)
 \end{aligned}$$

"WE'RE NOT TRYING TO MODEL ACTUAL PEOPLE,
 WE'RE MODELING RATIONAL PEOPLE."

11/5/18

ALL OF WHICH REDUCES TO:

$$x(1-\delta^T) = \left(\sum_{t=1}^T \delta^{t-1} x \right) \quad \delta \in [0,1]$$

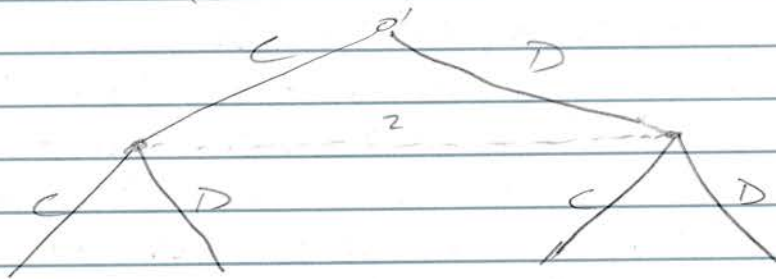
\swarrow 1 - PERFECT PATIENCE
 \searrow PERFECT 0-LIMIT

REDUCED FURTHER: $x(1-\delta^T) = \sum_{t=1}^T \delta^{t-1} x$

TOTAL PRESENT VALUE:
 (OF PAYOFF STREAM) $\frac{x}{1-\delta}$

(THIS IS GEOMETRIC DISCOUNTING, NOT HYPERBOLIC ($\beta\delta^T$))

TOTAL PRESENT VALUE OF THE GAMES ON PREV. PAGE
 $1 + \delta(0 \text{ or } -2)$



L P P

L $(4+\delta(0), 4+\delta(0)) (4+\delta(-3), 4+\delta(-1))$ (LEFT)

P $(4+\delta(-1), 4+\delta(-3)) (4+\delta(-2), 4+\delta(-2))$

x L P

L $(1+\delta(0), 1+\delta(0)) (1+\delta(-3), 1+\delta(-1))$ (RIGHT)

P $(1+\delta(-1), 1+\delta(-3)) (1+\delta(-2), 1+\delta(-2))$

FOR THIS GAME, AN EQUILIBRIUM STRATEGY TENDS TOWARDS COOPERATION.

$S_1 = D, P, P, P$
 $S_2 = D, P, P, P$



	Player 2: L, P		Player 2: L, P		Player 2: L, P		Player 2: L, P	
Player 1: L, P	L	P	L	P	L	P	L	P
L	$4 + \delta(0)$ $4 + \delta(0)$	$4 + \delta(-3)$ $4 + \delta(-1)$	$-1 + \delta(0)$ $5 + \delta(0)$	$-1 + \delta(-3)$ $5 + \delta(-1)$	$5 + \delta(0)$ $-1 + \delta(0)$	$5 + \delta(-3)$ $-1 + \delta(-1)$	$1 + \delta(0)$ $1 + \delta(0)$	$1 + \delta(-3)$ $1 + \delta(-1)$
P	$4 + \delta(-1)$ $4 + \delta(-3)$	$4 + \delta(-2)$ $4 + \delta(-2)$	$-1 + \delta(-1)$ $5 + \delta(-3)$	$-1 + \delta(-2)$ $5 + \delta(-2)$	$5 + \delta(-1)$ $1 + \delta(-3)$	$5 + \delta(-2)$ $-1 + \delta(-2)$	$1 + \delta(-1)$ $1 + \delta(-3)$	$1 + \delta(-2)$ $1 + \delta(-2)$

FIRST-STAGE PAYOFF + DISCOUNTED 2ND STAGE PAYOFF

↳ WHAT PLAYERS DO IN THE FIRST STAGE DOESN'T MATTER FOR THE SUBGAME.

THE REASONING THAT LED TO THE NE IN THE FIRST SUBGAME CAN BE APPLIED TO ALL 4.

WHAT ARE THE SUBGAME-PERFECT EQUILIBRIA THAT MOTIVATE COOPERATION?

$S_1 = \{C, L\}$ $S_2 = \{C, L\}$

A SINGLE PROFITABLE DEVIATION IS ENOUGH TO SHOW δ SUBGAME PERFECT.

STRATEGY PROFILES MUST NOW INCLUDE δ VALUES
 "IF δ IS $\geq \frac{1}{2}$, THEN..." (BOUNDARY CONDITION)

GAME THEORY 63 NOTES

1. SHOW WHAT NEEDS TO HAPPEN OFF THE PATH TO KEEP BOTH PLAYERS ON THE PATH.
2. COMPARE CONTINUATION VALUES (CETERIS PARIBUS)

PLAYERS WILL ALWAYS PLAY NASH EQUILIBRIA

"KNIFE-EDGE CONDITION" $\rightarrow \delta = 0$ ^{NO PATIENCE} OR $\delta = 1$ ^{MAXIMUM PATIENCE}

MULTIPLE NEs IN THE FINAL STAGE ARE REQUIRED FOR "HISTORICALLY-CONTINGENT" STRATEGIES - CREDIBLE THREATS REQUIRE MULTIPLE NEs.

↳ THIS IS WHY WE DIDN'T DO TWO PRISONERS' DILEMMAS.

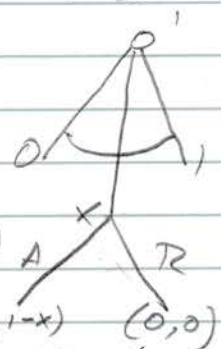
INFINITELY REPEATED P.D. RESULTS IN

- 1) MUTUAL COOPERATION; 2) INFINITE SUBGAME PERF. EQ.

DYNAMIC BARGAINING MODELS

ULTIMATUM GAME

PLAYER 1 OFFERS x
 PLAYER 2 ACCEPTS $(1-x)$
 OR REJECTS $(0,0)$



STRATEGY PROFILES

$S_1 =$ (ONE STRATEGY)

$S_2 =$ (INFINITE SUBGAMES)

P_2 IS PLAYING A CUT-POINT STRATEGY (MONOTONICALLY INCREASING IN x)

$P_1: BSR_1: \{ \text{PROPOSE } x^* \}$ $BSR_2: \{ \text{ACCEPT IF } x \geq x^* \}$
 REJECT OTHERWISE

$BSR_2(x) = \{ \text{ACCEPT IF } x < 1$
 ACCEPT OR REJECT IF $x = 1 \}$

$S_2': \text{ACCEPT } x < 1$
 REJECT $x = 1$

$S_2: \text{ACCEPT ALL } x$
 $S_1 = \text{PROPOSE } x = 1$

BUT IN THE CASE OF S_2' (ACCEPT $x < 1$
REJECT $x = 1$)

B_1 DOESN'T HAVE A BEST RESPONSE.

\hookrightarrow B/C $1 - \delta > 1 - \epsilon$ FOR ALL EPSILON.
SO NOT A SUBGAME PERFECT EQ.

NEW GAME:

(ADDITIONAL
STAGE)

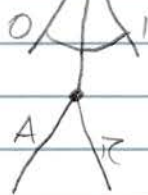


④ S_1 : AT $t=1$, PROPOSE $1 - \delta$

② AT $t=2$, ACCEPT ALL x

③ S_2 : @ $t=1$, ACCEPT IF $x \leq 1 - \delta$,
REJECT OTHERWISE

① AT $t=2$, PROPOSE $x=0$



WE CAN
SIMPLY
PUT
 $(0, \delta)$ HERE

FROM 1ST
STAGE $(\delta x, \delta(1-x))$

* SEE FIGURE 10 IN NOTES

NO EXAM QUESTIONS ON BARGAINING / INFINITE GAMES

GAME THEORY 7A

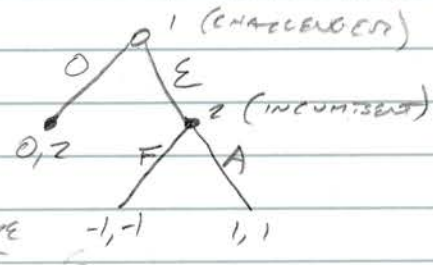
BAYESIAN GAMES

CANDIDATE ENTRY DETERMINANCE

TYPES

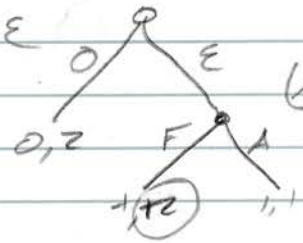
1 - CONTINGENT WAYS
2 TYPES → TIMID →

	1/2	F	A
0	0, 2	0, 2	0, 2
E	-1, -1	(1, 1)	



SUBGAME PERFECT

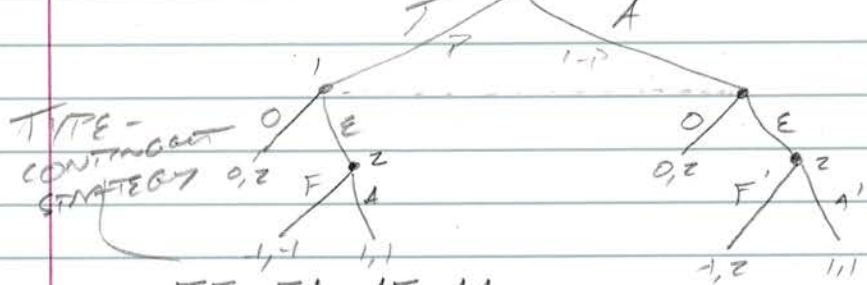
AGGRESSIVE



SAME EXCEPT FOR P2 PAYOFF IN F

TIMID: P
AGGRESSIVE: 1-P

DECIDES ON TYPE



TYPE-CONTINGENT STRATEGY

	FF	FA	AF	AA
0	0, 2	0, 2	0, 2	0, 2
E	$(-1, 2-3P)$	$(2P, 1-2P)$	$(2P-1, 2-P)$	(1, 1)

THESE 3 ARE BAYESIAN NASH-EQUILIBRIA

USE BACKWARDS INDUCTION BUT WEIGHT THE PROBABILITIES.

ASSOCIATE SINGLETON NODES W/ TYPES

TYPES REPRESENT UNCERTAINTY ABOUT EITHER (WHAT?)

Θ_1 = PLAYER 1'S TYPE SPACE = $\{\theta_1^1, \dots, \theta_1^k\}$

$P_i = P(\theta_i)$ = PLAYER I'S BELIEFS ABOUT OTHER

PLAYERS' TYPES = (θ_{-i})

Players will use BAYES' RULE TO UPDATE BELIEFS. (PRIOR BELIEFS BECOME POSTERIOR BELIEFS VIA AN UPDATING MECHANISM LIKE BAYES' RULE).

H = HYPOTHESES E = EVIDENCE

(T OR F)

(PERTINENT TO HYPOTHESIS)

CONDITIONAL PROBABILITY OF HYPOTHESIS GIVEN EVIDENCE

$$Pr(H|E) = \frac{Pr(H+E) = Pr(H) \times Pr(E|H)}{Pr(E) \quad Pr(H)Pr(E|H) + Pr(F)Pr(E|F)}$$

SPECIFY ACTIONS FOR EACH TYPE

$u_1(E, FA; \theta_2)$ (UTILITY OF PLAYING E AGAINST FA GIVEN PLAYER 2'S TYPE).

$$= \underbrace{p}_{\theta_2 = T}(-1) + \underbrace{(1-p)}_{\theta_2 = F}(1)$$

$$\phi_1(\theta_2 = T) \quad u_1(s_1 = E, s_2 = FA; \theta_2 = T)$$

GAME THEORY 8A

THE PAYOFF OF MULTIPLE TYPES WILL CHANGE ^{ACROSS} PAGES TABLE

	L	TR	L	TR
L	2,2	0,0	L	0,2
TR	0,0	4,4	TR	4,0

A WELL-FOUNDED STRATEGY TREE REQUIRES A SPECIFICATION FOR EVERY TYPE

IN THIS EXAMPLE, PLAYER 1 HAS 2 POSSIBLE TYPES (SO THEY HAVE 4 STRATEGIES)

TYPES: COOPERATIVE + CONTRARIAN (1-P)

	PLAYER 1		PLAYER 2	
A/B	L	TR	L	TR
LL	2P	2(1-P)	2	0
LTR	4-2P	0	2P	4-4P
TRC	0	2+2P	2-2P	4P
TRTR	4-4P	4P	0	4

WHEN THE VALUE OF P ISN'T SPECIFIED, WE CAN STILL FIND SOME BEST RESPONSES ($0 < P < 1$) AND WE CAN IDENTIFY CUT POINTS ($P > \frac{2}{3}$) WE CAN CREATE AN INEQUALITY FOR EACH ROW/col.

DRAW A NUMBER LINE: $0 \quad \frac{1}{3} \quad \frac{2}{3} \quad 1$

$$BR_2 = \begin{cases} L & \text{IF } LL, TR \\ R & \text{IF } LR, TR \end{cases}$$

IT'S POSSIBLE TO DO THIS ANALYSIS FOR $\text{TR} \geq 20$
 (BUT WE HAVE TO TREAT THE 4 OPTIONS
 USING PAIRWISE COMPARISONS)

FOR:

$P < 1$ LR IS TR TO L

$P = 1$ LL, LTR " "

$P \in (0, 1)$ TRL IS TR TO TR

$P = 1$ TRL, TRR

$P = 0$ LL, TRL

F = Firm

S.8 $\Theta_B = \{L, M, H\}$ $S_F = P \in [0, \infty)$ ($P = \text{PRICE}$
 OF LOANVIST)

$S_B(\Theta_B) : [0, \infty) \times \{L, M, H\} \rightarrow \{A, R\}$
 FOR INSTANCE, $(P=10, \Theta=M) \rightarrow \{R\}$
 (MUST SPECIFY FOR ALL P, Θ)

$$U_F(P, A; \Theta_B) = \begin{cases} 14 - P & \text{IF } L \\ 24 - P & \text{IF } M \\ 34 - P & \text{IF } H \end{cases}$$

$$U_F(P, R; \Theta_B) = 0$$

$$U_B(P, A; \Theta_B) = P \begin{cases} 10 & \text{IF } L \\ 20 & \text{IF } M \\ 30 & \text{IF } H \end{cases}$$

$$U_B(P, R; \Theta_B) = \begin{cases} 10 & \text{IF } L \\ 20 & \text{IF } M \\ 30 & \text{IF } H \end{cases}$$

WHEN IN DOUBT, START FIGURING OUT BEST RESPONSES

$$TR(\Theta_B) = \begin{cases} A & \text{IFF } P \geq 10 + \Theta_B = L \\ A & \text{IFF } P \geq 20 + \Theta_B = M \\ A & \text{IFF } P \geq 30 + \Theta_B = H \end{cases}$$

$$EV(\text{ACCEPTED } T) = \frac{1}{3}(14 + 24 + 34) = 24$$

$$EV(\text{RESERVE}) = \frac{1}{3}(10 + 20 + 30) = 20$$

BUT THE HIGH TYPE WOULDN'T ACCEPT 20

$$\text{NEW } EV(\text{ACCEPTED } P) = \frac{1}{2}(14 + 24) + 0(34) = 19$$

GAME THEORY 8B

IF YOU PURSUE THIS UPDATING LOGIC "AC TTS WAY DOWN",
THE FINAL BARGAINING RANGE IS 10-14.

$$TSNE = S_B^*(\Theta) = \begin{cases} A \text{ iff } P \geq P^* \in [10, 14] & \text{if L} \\ A \text{ iff } P \geq 20 & \text{if M} \\ A \text{ iff } P \geq 30 & \text{if H} \end{cases}$$

$$S_F^* = P^* \in [10, 14]$$

(INFINITE TSNE, BTW 10+14)

INFINITE TYPES CASE: TYPES DISTRIBUTED

$$U_F(P, S_B; \Theta) = \begin{cases} 0 & \text{if } L \\ 30 - P & \text{if } \{A\} \\ P & \text{if } A \end{cases}$$

$$U_B(P, S_B; \Theta) = \begin{cases} P & \text{if } A \\ 20 & \text{if } R \end{cases}$$

$P \geq 20$ LOWER BOUND

$$E[S_B(P)] = \frac{1}{2}P \text{ (THRESHOLD TYPE)} \quad T = \frac{1}{2}$$

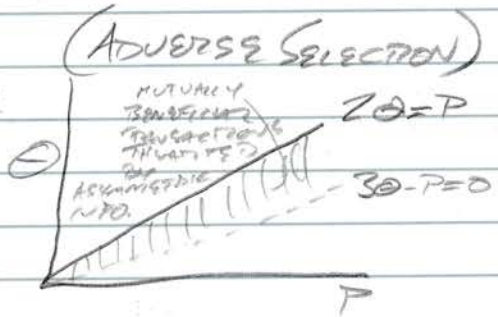
$$0 \leq P \leq (3 \times E[V[\Theta]] - P) \text{ (MIDPOINT OF INTERVAL)}$$

$$0 \leq P \leq (3 - P) \rightarrow \text{UPPER BOUND}$$

$$\rightarrow = \begin{cases} 1 & \text{if } \Theta \leq \frac{1}{2}P \\ 0 & \text{if } \Theta > \frac{1}{2}P \end{cases} \quad \downarrow \quad 3\left(\frac{P}{4}\right) - P = \frac{1}{4}P$$

$$TSNE = S_F^* = P = 0$$

$$S_B^* = \begin{cases} A & \text{if } \Theta \leq \frac{1}{2}P \\ R & \text{otherwise} \end{cases}$$



ASYMMETRIC INFORMATION HAS
BIG RAMIFICATIONS FOR WHAT
IS RATIONAL.

WRITE TOSS
 NEW POWERS
 ALLOWING TOSS
 WINGS = SAM
 RESPOND TO
 TIME + GAT

$$1 - (1-u)\sqrt{c}^{1/2}$$

PUBLIC GOODS CONTRIBUTION GAME

$e_i \in \{0, 1\}$ $e_2 \in \{0, 1\}$ (TWO LEGISLATORS)
 $c \in (0, 1]$

$$u_i(e_i, e_j, \theta_i) = \begin{cases} \theta_i^2 - c & \text{if } e_i = 1 \\ \theta_i^2 & \text{if } e_i = 0 + e_j = 1 \\ 0 & \text{if } e_i = e_j = 0 \end{cases}$$

FIRST, FIGURE OUT BEST RESPONSES:

$$\theta_i^2 - c \geq \text{PR}(e_2 = 1 | \theta_2) \theta_i^2$$

$$(1 - \text{PR}(e_2 = 1 | \theta_2)) 0 - \text{CANCELS}$$

SYMMETRIC FOR P2

$$\theta_i \geq \sqrt{\frac{c}{1 - \text{PR}(e_2 = 1 | \theta_2)}} \quad (\text{THRESHOLD TYPE})$$

$$\hat{\theta}_1 \equiv \sqrt{\frac{c}{1 - (1 - \hat{\theta}_2)}} = \sqrt{\frac{c}{\hat{\theta}_2}}$$

$$\hat{\theta}_1 \equiv \sqrt{\frac{c}{\hat{\theta}_2}} \quad \hat{\theta}_2 \equiv \sqrt{\frac{c}{\hat{\theta}_1}}$$

So: $\hat{\theta}_2^2 \hat{\theta}_1 = c$
 And: $\hat{\theta}_1 = \hat{\theta}_2$

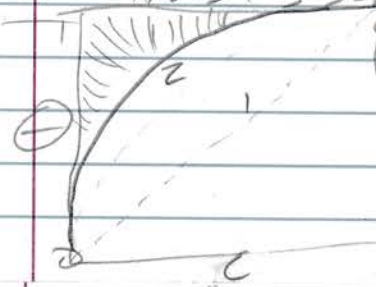
IF THE PROBLEM IS FULLY SYMMETRIC, THEN $\hat{\theta}_1 = \hat{\theta}_2$

GOAL: EXPRESS $\hat{\theta}$ IN TERMS OF c

$$\hat{\theta}_1(\hat{\theta}_2) = \begin{cases} c^{1/3} & \text{if } \hat{\theta}_2 \geq c \\ 1 & \text{if } \hat{\theta}_2 < c \end{cases} \quad \text{TRUE: IF } \hat{\theta}_1 \geq c$$

$$s_i^*(\theta) = \begin{cases} e_i = 1 & \theta_i \geq \hat{\theta}_i \\ e_i = 0 & \theta_i < \hat{\theta}_i \end{cases} \quad \text{FOR ALL } i, \text{ THEN } s_i^*(\theta) = \begin{cases} e_i = 1 & \text{if } \theta \geq c^{1/3} \\ e_i = 0 & \text{otherwise} \end{cases}$$

3-PLAYERS: 3 SCENARIOS FOR FREE-RIDING



$E[U_i]$ IS NEGATIVE ON ... (1) $\theta^2 - c = 0$

BREAK-EVEN TYPE: (2) $\theta^2 = c$

PROBABILITY OF PROVISION: $\hat{\theta} = c^{1/3}$

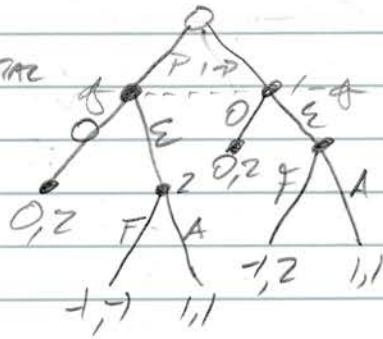
$$\pi = (1 - c^{1/3})^2 + 2(1 - c^{1/3})c^{1/3}$$

$$= 1 - c^{2/3}$$

GAME THEORY - 9A

11 / 27 / 18

WE USE SEQUENTIAL RATIONALITY TO ELIMINATE $\frac{2}{3}$ OF THE NASH EQUILIBRIA

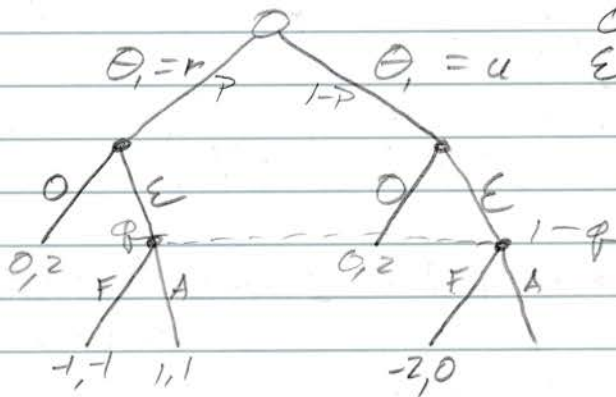


SEQ. RATIONAL BNE ($P = \frac{2}{3}$)

- x (O, F)
- x (O, FA)
- ✓ (E, AF)

BEST RESPONSE (WITH TYPES) MUST TAKE INTO ACCOUNT UNCERTAINTY ABOUT THE OTHER PLAYERS' TYPES ($q, 1-q$).

Now P_1 HAS THE TYPES: $P = \frac{1}{2}$ FF FA AF AA



OOFF IS SUBGAME PERFECT BUT NOT SEQ. RATIONAL

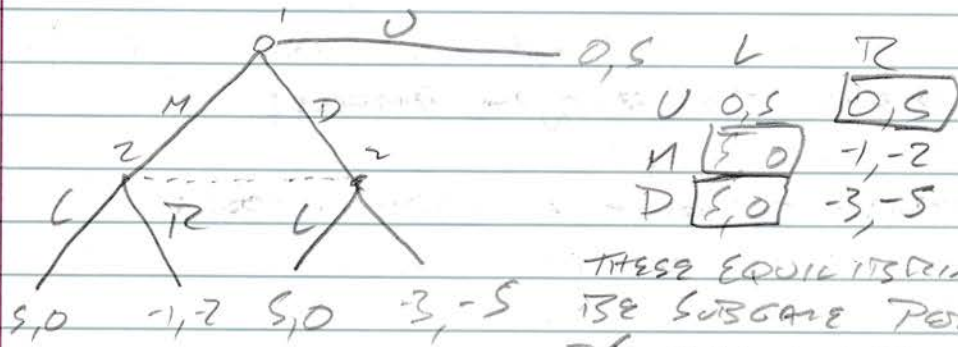
$P = \frac{1}{2}$

	F	A
OO	0, 2	0, 2
OE	-1, 1	$\frac{1}{2}, \frac{3}{2}$
EO	$-\frac{1}{2}, \frac{1}{2}$	$\frac{1}{2}, \frac{3}{2}$
EE	$-\frac{3}{2}, -\frac{1}{2}$	1, 1

2 BNE: (OO, F) + (EE, A)

THERE'S NO BELIEF FOR WHICH FIGHT IS A BEST RESPONSE.

NEW CONCEPT OF RATIONALITY: PLAYERS ARE PLAYING BEST RESPONSES AT EACH OF THEIR INFORMATION SETS GIVEN THEIR BELIEFS RE: OTHER PLAYERS' STRATEGIES.



THESE EQUILIBRIA MUST BE SUBGAME PERFECT B/C ONLY ONE PROPER EQUILIBRIUM

"SYSTEM OF BELIEFS" \rightarrow A WELL-FOUNDED PROBABILITY DISTRIBUTION ASSIGNED TO EACH INFORMATION SET.



$$q = \Pr(\theta = t | \sigma, (E|r))$$

$$1-q = \Pr(\theta = b | \sigma, (E|r))$$

$$q = \Pr(\theta = t | E)$$

$$1-q = \Pr(\theta = ?)$$

$$q = \Pr(\theta = t | E) = \frac{\Pr(\theta = t) \Pr(E|r)}{\Pr(E)}$$

$$= \frac{\frac{1}{2} \times 1}{\frac{1}{2}} = 1$$

LEARNING ONLY TAKES PLACE (BELIEFS ONLY CHANGE FROM THEIR INITIAL VALUES) WHEN

PERFECT BAYESIAN EQUILIBRIUM (σ^* , μ)
 APPLIES TO STRATEGY PROFILES

CHARACTERIZING A PBE INVOLVES STATING
 THE STRATEGY REQUIREMENTS AND STATING
 WHAT THE BELIEFS ARE.

PROBS. DIST'D
 ASSIGNED TO EACH
 INFORMATION SETS.

REQUIREMENTS:

- #1 - μ IS A WELL-DEFINED SYSTEM OF BELIEFS
- #2 - μ IS CONSISTENT W/ BAYES' RULE AT ALL INFORMATION SETS ON THE PATH OF σ^* .
- #3 - μ IS CONSISTENT WITH BAYES' RULE AT ALL OFF-PATH INFORMATION SETS WHERE POSSIBLE
- #4 - σ_i^* MUST BE SEQUENTIALLY RATIONAL AT EACH INFORMATION SET (ON OR OFF THE PATH) GIVEN $\sigma_{-i}^* + \mu_i$.

(σ^* , μ) IS A PBE IFF #S 1-4 OBTAIN.

WEAK PBE ELIMINATES #3 - FOR INFORMATION SETS OFF THE PATH, ANYTHING GOES.

(EXERC. PAGE)

EXAMPLES: $u_1(F, \varphi) \geq u_1(A, \varphi)$ (NEVER TRUE)

BACKWARDS
 INDUCTION
 ↓

$BR_2(\varphi) = \{A\}$ FOR ALL VALUES OF φ

$BR_1(s_2; \theta_1) = \{E\}$ IF $\theta_1 = r$

$\{E\}$ IF $\theta_1 = u$

$= \{E\}$ FOR ALL θ_1

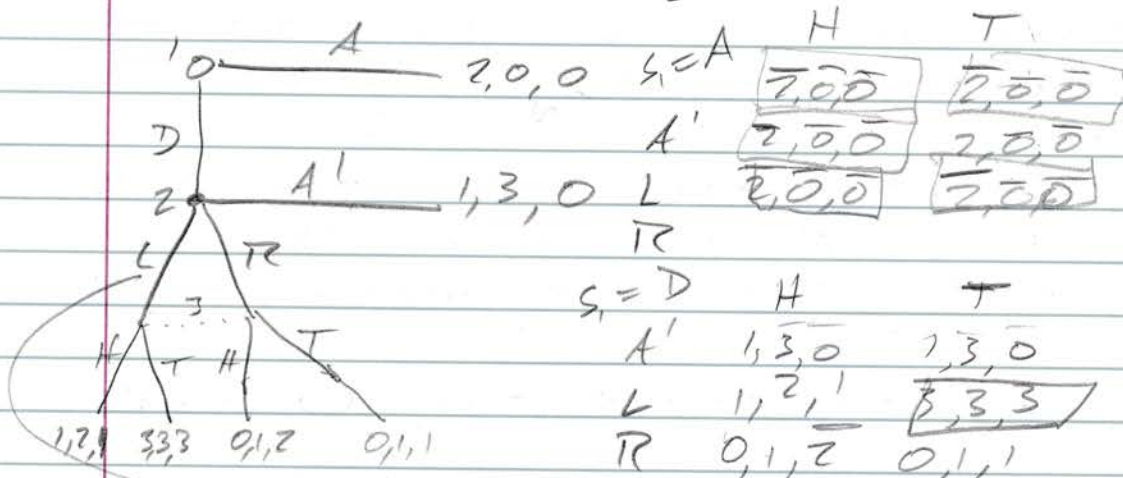
$$\sigma^* \begin{cases} s_1(\theta_1) = E \text{ FOR ALL } \theta_1 \\ s_2 = A \end{cases}$$

Now we need beliefs:

Since $S_1(0,1)$, $\mu = \phi = P$ (NO UPDATING
 FROM PEERS: TYPES ARE EQUALLY LIKELY).

So $(0,0,F)$ CAN BE ELIMINATED (NOT PBE)
 B/C IT DOESN'T CONFORM TO σ^*

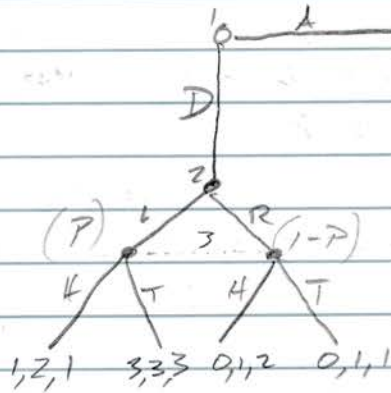
REQUIREMENT #3 IS TRICKY - SEE EX.
 4.0.2 IN THE NOTES.



L STRICTLY DOMINATES R, SO PLAYER 2
 WILL NEVER PLAY R. PLAYER 3 SHOULD KNOW THIS.

THE EXAMPLE ABOVE IS WRONG - HE TR- DID
 IT ON 1/2S.

GAME THEORY 9B



- (IS) NE: (A, L, H)
 (A, R, H)
 (A, R, T)
 (D, L, T)

SUBGAME PERFECT: (D, L, T)
 PERFECT (IN EACH SUBGAME)

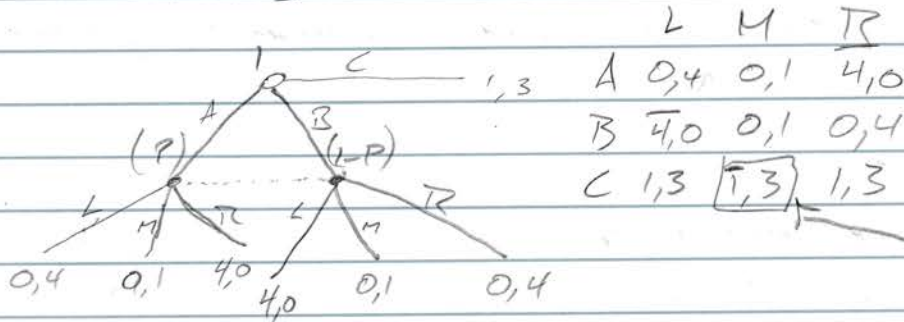
OFF-THE-PATH STRATEGIES, ONCE UPDATED BY BAYES' RULE, CONSTRAIN POSSIBLE EQUILIBRIA.

PSE: $(D, L, T) + (P=1)$
 (MUST HAVE BOTH PARTS)

$$BSR_3(P) = \begin{cases} \{H\} & \text{if } P < \frac{1}{3} \\ \{H, T\} & \text{if } P = \frac{1}{3} \\ \{T\} & \text{if } P > \frac{1}{3} \end{cases}$$
 POSSIBILITIES FOR T MAY BE CONSTRAINED BY OFF-THE-PATH STRATEGIES

* PROCEDURE: FIND THE SPE, THEN ASK IF THERE ARE ANY BELIEFS THAT JUSTIFY THOSE STRATEGIES.

* [SEE 4.0.2] - STATE BSR_3 AS A FUNCTION OF BELIEFS



	L	M	T
A	0,4	0,1	4,0
B	4,0	0,1	0,4
C	1,3	1,3	1,3

IF WE TREAT P AS A LOTTERY BETWEEN L + TR, THEN THE AVERAGE OF THAT MIX WILL BE BETTER THAN M.

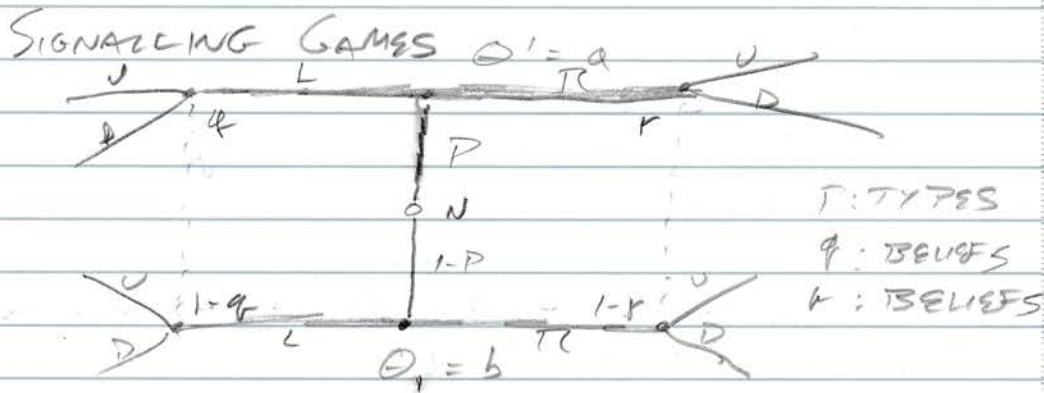
$$BSR_3 \begin{matrix} L \\ T \end{matrix} \quad \begin{matrix} P(L) \geq 1 \\ P \geq 1/4 \end{matrix} \quad \begin{matrix} 4 - 4P \geq 1 \\ 3 \geq 4P \end{matrix} \Rightarrow \begin{matrix} \text{THIS RULES OUT} \\ \text{THE NE AT C, M} \end{matrix}$$

$$\frac{1}{4} \quad \frac{3}{4}$$

$$\frac{3}{4} \geq P$$

So we try a mixed strategy: $r = \frac{1}{2}$; $q = \frac{1}{2}$
 (because the top-left 4 cells = matching pennies)
 \hookrightarrow Results in $EU_{1+2} = 2$

So we have a new PBE: $\sigma_1 = r = \frac{1}{2}$; $\sigma_2 = q = \frac{1}{2}$



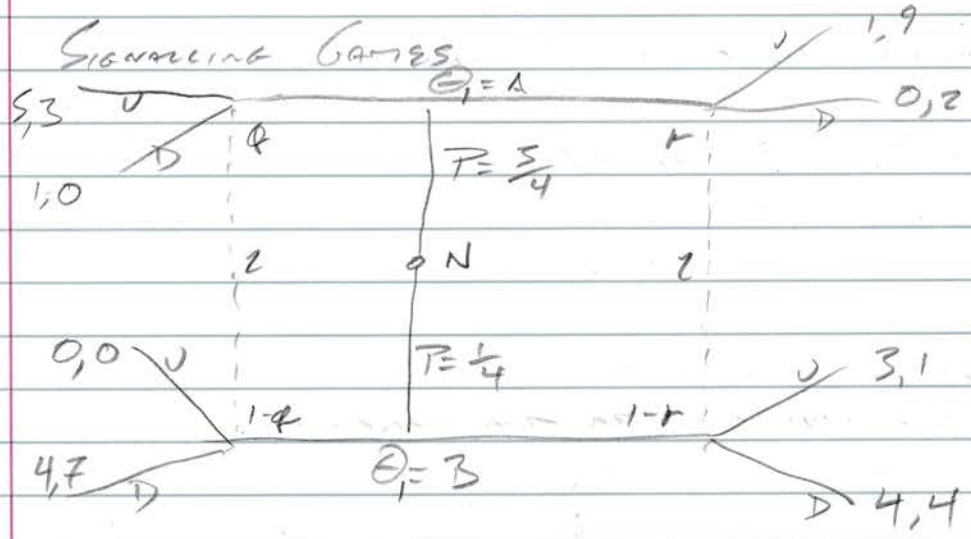
PLAYERS SEE THE OTHER PLAYER'S ACTIONS,
 BUT THEY DON'T KNOW WHICH TYPE PERFORMED IT.

SEPARATING STRATEGY - TYPES PERFORM DIFFERENT ACTIONS.
 IF THIS RESULTS IN AN EQUILIBRIUM, IT IS A SEPARATING EQUILIBRIUM
 OR A POOLING EQ

PERFECTLY INFORMATIVE / OR UNINFORMATIVE POOLING EQUILIBRIUM -
 SEPARATING EQUILIBRIUM -

IMPERFECTLY INFORMATIVE SEMI-SEPARATING EQUILIBRIUM -

GAME THEORY 10A



Look for Player 2's Best Response as a function of their information set.

#1 $U_2(U, p) \geq U_2(D, p)$
 $3p + 0(1-p) \geq 0p + 7(1-p)$
 $3p \geq 7 - 7p$
 $10p \geq 7; p \geq \frac{7}{10}$

#2 $U_2(U, r) \geq U_2(D, r)$
 $4r + (1-r) \geq 2r + 4(1-r)$
 $8r + 1 \geq 4 - 2r$
 $10r \geq 3; r \geq \frac{3}{10}$

So $BR_1 =$
 $BR_2(p) = \begin{cases} \{U\} & p \geq \frac{7}{10} \\ \{U, D\} & p = \frac{7}{10} \\ \{D\} & p < \frac{7}{10} \end{cases}$
 $BR_2(r) = \begin{cases} \{U\} & r \geq \frac{3}{10} \\ \{U, D\} & r = \frac{3}{10} \\ \{D\} & r < \frac{3}{10} \end{cases}$

Look for Pooling Equilibria - Both types take the same action. Two in this case: (Give no information on types)

$q = Pr(A|L) = \frac{Pr(A)Pr(L|A)}{Pr(L)} = q = \frac{3}{4}, r = \text{undef.}$

So: Against $\{L\}$, P_2 plays $\{U\}$ ($\sigma(A)=L, B=1$)

Now we look for the strategy that Player 2 would have to play off the path to sustain the pooling equilibrium.

So LL is not a pooling equilibrium

Now we check for a second pooling equilibrium - at TR, TR ($\theta = A$ plays R)
 $r = \frac{3}{4}$, $q = \text{UNDEF.}$ ($\theta = B$ plays R)
 we can rule out U ($s > 1$) and d ($z > 1$)
 so TR, TR is not a pooling equilibrium.

Now we check for separating equilibria.
 There will also be two here: L, R and R, L
 First we check L, R .

Since we know θ_A will play L , then $q = 1$
 $r = 0$

So against L , $\sigma_1 = u$; against R , $\sigma_1 = d$
 which means this is SSE as long as there
 are no profitable deviations. (there is one?)

Now we check R, L : $q = 0$, $r = 1$

So against L , $\sigma_1 = d$; against R , $\sigma_1 = u$
 neither player

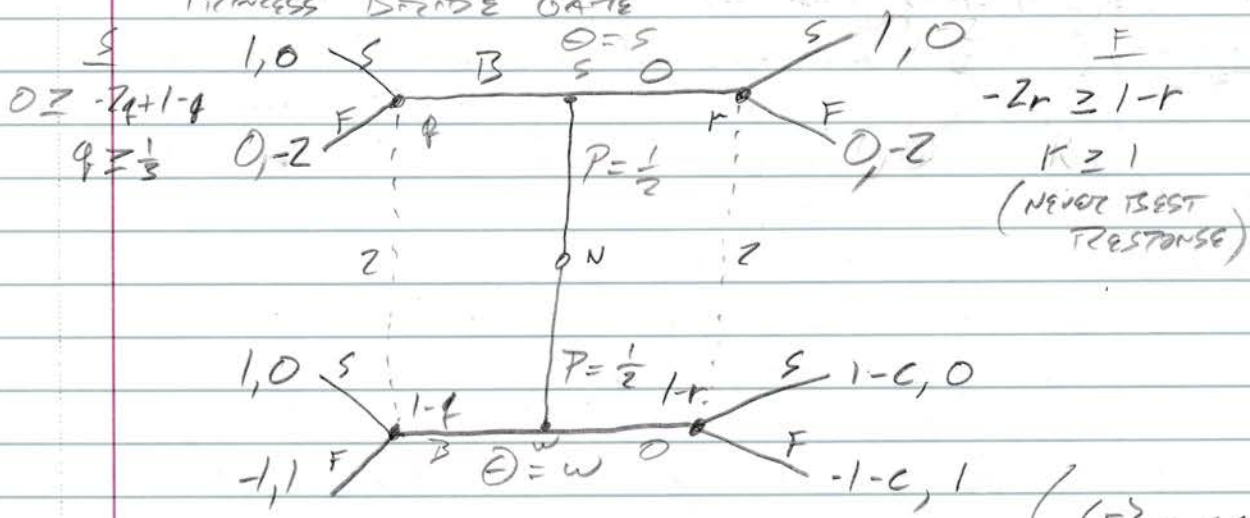
Semi-separating: $q = Pr(A|L) = \frac{Pr(A)Pr(L|A)}{Pr(L)}$
 $(Pr(L) = Pr(A)Pr(L|A) + Pr(B)Pr(L|B))$

Solve this, and you get an answer with s in it
 (s is the mixed strategy played by one of π 's types)
 \rightarrow then set this equal to the q -value ($\frac{3}{4}$)

q is a function of the probability of playing L .

GAME THEORY 1013

"PRINCESS BRIDE GAME"



$BTR_2(\emptyset) = \{F\}$ For $\forall r$; $BTR_2(r) =$

AGAINST \emptyset , $s_2 = S$; AGAINST B , $s_2 = F$

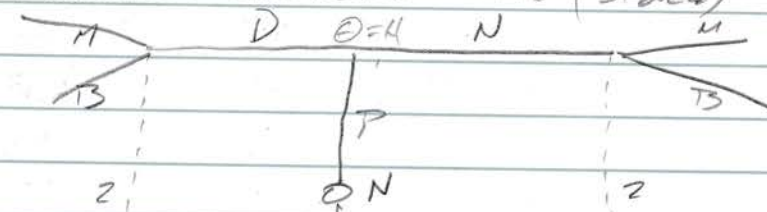
$\{F\}$ if $q < \frac{1}{3}$
 $\{F, S\}$ if $q = \frac{1}{3}$
 $\{S\}$ if $q > \frac{1}{3}$

$1 - c \geq -1$ } TYPE W HAS NO INCENTIVE TO DEVIATE
 $2 \geq c$ } AS LONG AS $c \leq 2$. THIS SUSTAINS THE EQ.

C VALUES MUST BE SPECIFIED AS PART OF A TBE
 "IF C IS SUFFICIENTLY HIGH, THEY SEPARATE. IF C IS SUFFICIENTLY LOW, THEY POOL."

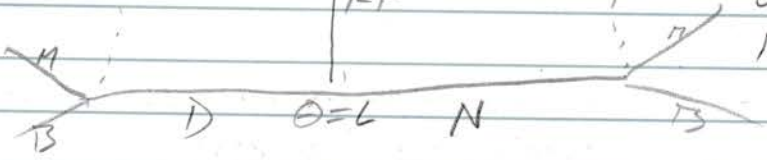
EDUCATION SIGNALING GAME (SPACE)

$W_H - C_H$
 $20 - W_H$
 $W_B - C_H$
 $11 - W_B$



W_M
 $20 - W_M$
 W_B
 $11 - W_B$

$W_H - C_L$
 $10 - W_H$
 $W_B - C_L$
 $9 - W_B$



W_M
 $10 - W_M$
 W_B
 $9 - W_B$

PLAYER 1'S TYPE AFFECTS BOTH PLAYERS' PAYOFFS.
 $P_1 =$ EMPLOYEE, $P_2 =$ MANAGER

ASSUMPTIONS: $C_L > C_H > 0$; $W_M > W_B$

SPECIFY THE EQUILIBRIUM WE WANT TO SOLVE FOR
 BEFORE WORKING OUT THE BEST-RESPONSE FUNCTION
 (OTHERWISE WE'LL HAVE TWO UNKNOWN PARAMETERS)

SO WE CONSTRUCT A STRATEGY FOR PLAYER 1:

$\sigma_1(H) = D$ so $q+r$ are: $q=1; r=0$
 $\sigma_1(L) = N$ so $20 - W_H \geq 11 - W_B$; $9 \geq W_H - W_B$

$BR_2(q=1) = \begin{cases} H & \text{if } 9 \geq W_M - W_B \\ B & \text{if } 9 < W_M - W_B \end{cases}$

$10 - W_M \leq 9 - W_M$; $1 \leq W_M - W_B$ | $9 \geq W_M - W_B \geq 1$
 $W_H - C_H \geq W_B$; $W_H - W_B \geq C_H$ (SUSTAINS SEPARATION)

SO: $P_2: 9 \geq W_M - W_B \geq 1$
 $P_1: C_L \geq W_M - W_B \geq C_H$

GRAPH THE PARAMETER SPACE

$BR_2 = \begin{cases} \text{AGAINST D, PLAY H} \\ \text{AGAINST N, PLAY B} \end{cases}$

IF $9 \geq C_L \geq W_M - W_B \geq C_H \geq 1$, THEN THIS EQ. IS TRUE